

Graph Theory Lecture 15

Dec. 26 -

Theorem 21 (Equitable edge-coloring)

If G has a k -edge-coloring, then G has an equitable edge-coloring, i.e., for any two $i, j \in \{1, 2, \dots, k\}$, $||f^{-1}(i)| - |f^{-1}(j)|| \leq 1$.

Proof. If there exist i and j such that $|f^{-1}(i)| - |f^{-1}(j)| \geq 2$, then

we consider the graph H induced by the set of edges colored i and j .

Then, H is a subgraph of G such that each component of H is

either a path or an even cycle. Since i occurs more times than

j , there exists an i - j path: $\overset{i}{\circ} - \overset{j}{\circ} - \overset{i}{\circ} - \overset{j}{\circ} - \dots - \overset{i}{\circ}$ whose

end edges are colored i . Now, by switching the colors on this path,

we obtain a new edge coloring of G such that i occurs one less

time and j occurs one more. It turns out that we can obtain

a k -edge-coloring s.t. $||f^{-1}(i)| - |f^{-1}(j)|| \leq 1$. As a consequence,

we are able to adjust all of them and obtain an equitable

k -edge-coloring. ■

(*) This theorem is also not difficult to prove, but very useful.

(o) Without using 4CT, the proof of Theorem 70 is very difficult.

(*) It was conjecture that if G is planar and $\Delta(G)$ is large enough, say, _____.

Theorem 22

The complete graph K_n is of Class 2 if and only if K_n is overfull or equivalently n is odd.

Proof. First, we claim that for each $m \geq 1$, K_{2m} is of Class 1. It suffices to give a $(2m-1)$ -edge-coloring of K_{2m} . For convenience, let

$V(K_{2m}) = \mathbb{Z}_{2m} = \{0, 1, 2, \dots, 2m-1\}$. For each color $i \in \{1, 2, \dots, 2m-1\}$,

let the set of edges colored i be

$$F_i = \{ (0, i), (i+1, i-1), (i+2, i-2), \dots, (i+m-1, i-m+1) \} \pmod{2m-1}.$$

See Figure 46 for an example of $m=5$ and $i=3$.

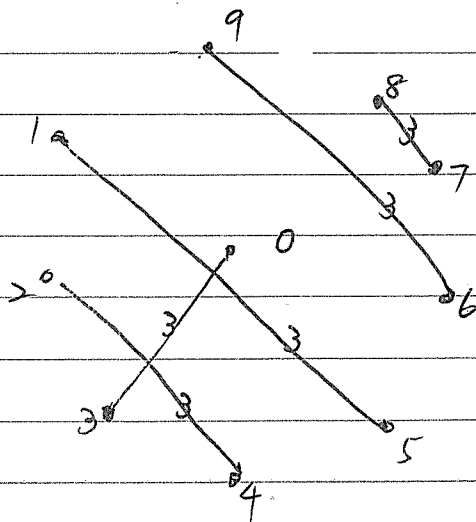


Figure 46. $\chi'(K_{10})=9$.

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Since $\Delta(K_{2m}) = 2m-1$, $\chi'(K_{2m}) = 2m-1$.

Now, by deleting v in K_{2m} , we obtain a $(2m-1)$ -edge-coloring of K_{2m-1} . On the other hand, it is not difficult to check that K_{2m-1} is overfull for $m \geq 2$, this concludes that $\chi'(K_{2m-1}) > \Delta(K_{2m-1}) = 2m-2$. ■

(\circ) This theorem is not difficult to prove, but it is very useful in the construction of "Combinatorial Designs".

($\circ\circ$) Equivalently, K_{2m} can be decomposed into $2m-1$ 1-factors, which is also known as a 1-factorization of K_{2m} .

($\circ\circ\circ$) If G is an r -regular graph and $\chi'(G) = r$, then G has a 1-factorization.

($\circ\circ\circ\circ$) It was conjectured that if G is r -regular and $r \geq \frac{|G|}{2}$, then G has a 1-factorization or equivalently $\chi'(G) = r$.

Theorem 21' (D. Hoffman et al.)

A complete multipartite graph G is of Class 2 if and only if G is overfull.

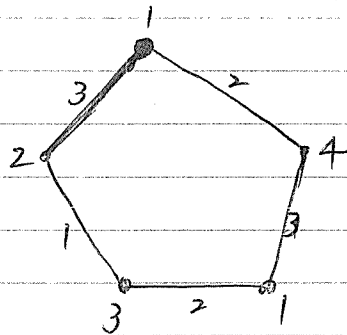
(*) Total coloring

A k -total coloring of a graph G is a mapping

$$\psi: V(G) \cup E(G) \rightarrow \{1, 2, \dots, k\} \text{ such that}$$

- (i) adjacent vertices receive distinct images,
- (ii) incident edges receive distinct images, and
- (iii) each vertex has a distinct image with its incident edges.

e.g.



A 4-total coloring of C_5 .

$$(*) \chi''(G) = \min. \{k \mid G \text{ has a } k\text{-total coloring}\}.$$

(Total chromatic number)

Theorem 23

$$\chi''(K_{2n+1}) = \chi''(K_{2n}) = 2n+1.$$

$$\chi''(K_5) = 5 \quad (?)$$

Proof. $\chi''(K_{2n+1})$ can be obtained by using $\chi'(K_{2n+1}) = 2n+1$.

TCC Conjecture $\chi''(G) \leq \Delta(G) + 2$.

Note that $\chi''(G) \geq \Delta(G) + 1$. As to the total coloring of K_{2n} , we claim that $2n$ colors are not enough.

Observe that each color class has at most one vertex and $n-1$ edges. So, $2n$ color classes will contain at most $2n$ vertices and $2n(n-1)$ edges. Hence, ^{there are} $2n^2$ elements (vertices and edges) in total. But, K_{2n} has $2n + \frac{2n(2n-1)}{2}$ elements to color, which is $2n^2 + n$. Clearly, $2n$ color is not enough. Since K_{2n+1} is $(2n+1)$ -total colorable, K_{2n} is also $(2n+1)$ -total colorable. The proof follows. \blacksquare

(c) Based on TCC Conjecture, a graph G is called Type 1 if $\chi''(G) = \Delta(G) + 1$ and Type 2 otherwise.

Theorem 24. $K_{m,n}$ is of Type 1 if and only if $m \neq n$.

Proof. (\Rightarrow) If $m = n$, then there are $2n + n^2$ elements to color. Since each color class contains at most n elements, $\Delta(G) + 1 = n + 1$ colors are not enough, $n(n+1) < 2n + n^2$. Hence, $\chi''(K_{n,n}) \geq n + 2$, and thus $K_{n,n}$ is ^{not} of Type 1.

(\Leftarrow)

On the other direction, let $m = n+k$. Now, $\Delta(K_{m,n}) = n+k$.

By the edge-coloring of $K_{m,n}$, we have an $\sqrt{n \times (n+k)}$ Latin rectangle based on $\{1, 2, \dots, n+k\}$, see Figure 47. Since $k \geq 1$, we may extend this rectangle to $(n+1) \times (n+k)$ and the last row can be used to color the vertices of

A. Finally, color all vertices of B by one extra color, we have

$$\chi''(K_{m,n}) \leq n+k+1 = \Delta(K_{m,n})+1.$$

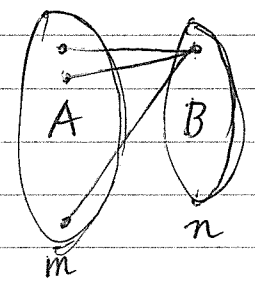
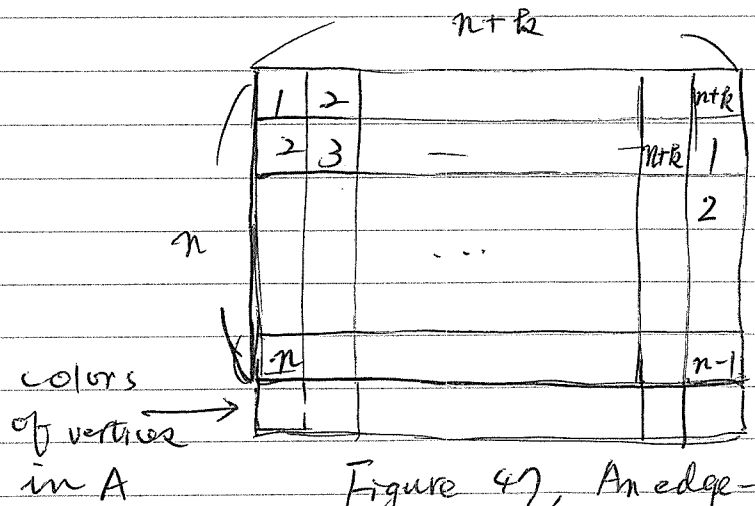
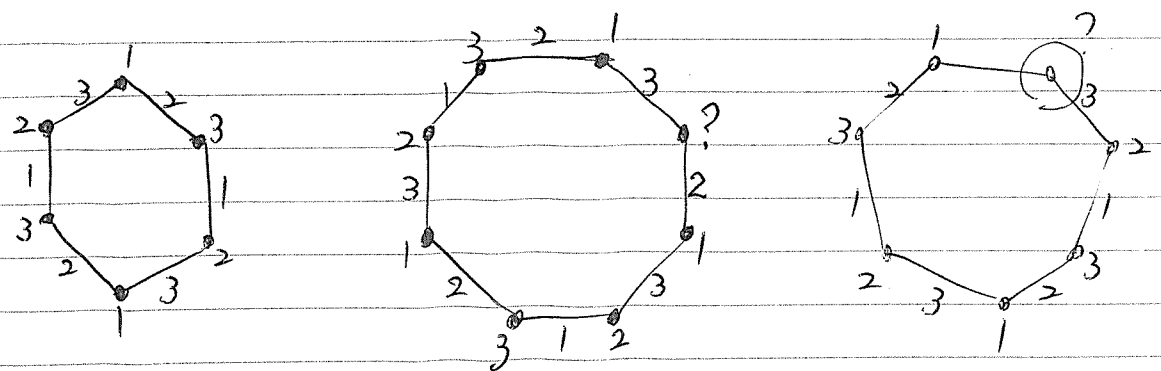


Figure 47, An edge-coloring of $K_{m,n}$.

(*) A cycle C_n is of Type 1 if and only if $n \equiv 0 \pmod{6}$.



If we use three colors, then starting from one vertex and one edge, all the colors of the others are forced!

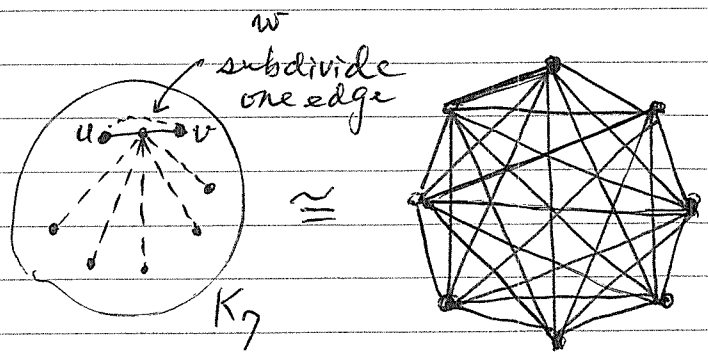
- (*) The deficiency of a graph G , $def(G)$, is defined as $\sum_{v \in V(G)} (\Delta(G) - \deg(v))$.
- (**) G is conformable if G has a vertex coloring $\varphi : V(G) \rightarrow \{1, 2, \dots, \frac{\Delta(G)}{2} + 1\}$ such that $def(G) \geq |\{i \mid |G| - |\varphi^{-1}(i)| \equiv 1 \pmod{2}\}|$.

Theorem 75

Let S_i be a star with i edges. Then, $K_{2n} - S_1 - S_{2n-3}$ is of Type 2.

$\parallel G(1, 2n-3)$

Proof. Assume that $G(1, 2n-3)$ is of Type 1, i.e., there exists a total coloring ψ of $G(1, 2n-3)$ using $\Delta(G)+1 = 2n-1$ colors. Let uv be the edge subdivided, see Figure 48. First, if $\psi(u) = \psi(v)$, then let $\psi(w) = 1$, $\psi(uw) = 2$ and $\psi(wv) = 3$. Let r_j the number of vertices in which j occurs in either v or an edge incident to v . Hence, $\sum_{j=0}^{2n-2} r_j = |K_{2n}| + 2 \parallel G(1, 2n-3) \parallel$. (?) Now, 0 occurs in at most $2n-2$ vertices. Hence, $\sum_{j=0}^{2n-2} r_j \leq (2n-2) + 3 \cdot (2n) + (2n-5) \cdot (2n-1) = 4n^2 - 4n + 3 < |K_{2n}| + 2 \parallel G(1, 2n-3) \parallel$.



$n = 4$
 $\leftarrow G(1, 5)$
 Figure 48. $G(1, 5)$.

On the other hand, if $\psi(u) \neq \psi(v)$, a similar argument shows that $2n-1$ colors are not enough. Hence, $G(1, 2n-3)$ is of Type 2. ■

(*) If G is of Type 1, then G is conformable.

Proof. If there exists a color i such that $|G| - |\psi_{V(G)}^{-1}(i)| \equiv 1 \pmod{2}$, then i occurs in at most $|G|-1$ vertices. Since every color occurs around a major vertex, $\Delta(G)+1$ colors are not enough if $\text{def}(G) < |\{i \mid |G| - |\psi_{V(G)}^{-1}(i)| \equiv 1 \pmod{2}\}|$. So, if G is not conformable, then G is not of Type 1. ■

Conjecture (Chetwynd and Hilton, 1988)

Let G be a simple graph with $\Delta(G) \geq \lfloor \frac{|G|+1}{2} \rfloor$. Then, G is of Type 2 if and only if there exists a non-conformable subgraph H of G such that $\Delta(H) = \Delta(G)$ (and G is not the Chen and Fu graph). ↑ 後者不可

(000) The ^Vconjecture was disproved by using Theorem 75. (original)

(*) The graph $G(1, 2n-3)$ is known as Chen and Fu graph.