

(*) $|x-y| \leq 1$ and $x+y \leq 2^k$, let $x \leq y$.

$$\Rightarrow y-x \leq 1 \text{ and } x+y \leq 2^k$$

$$\Rightarrow x \geq y-1 \text{ thus } 2y-1 \leq 2^k, 2y \leq 2^k+1.$$

If $y > 2^{k-1}$, then $2y \geq 2^k+2$. Hence $y \leq 2^{k-1}$, so is x .

Lemma Let (a, b) be a state such that $b \geq a-1 \geq 1$. Then there exists a test in this state yielding states (a_1, b_1) and

(a_2, b_2) such that

$$(1) \lfloor \frac{a}{2} \rfloor \leq a_1 \leq \lfloor \frac{a+1}{2} \rfloor, \lfloor \frac{a}{2} \rfloor \leq a_2 \leq \lfloor \frac{a+1}{2} \rfloor;$$

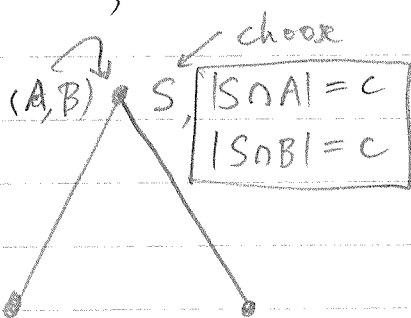
$$(2) b_1 \geq a_1-1, b_2 \geq a_2-1; \text{ and}$$

$$(3) \text{ch}(a_1, b_1), \text{ch}(a_2, b_2) \leq \text{ch}(a, b) - 1.$$

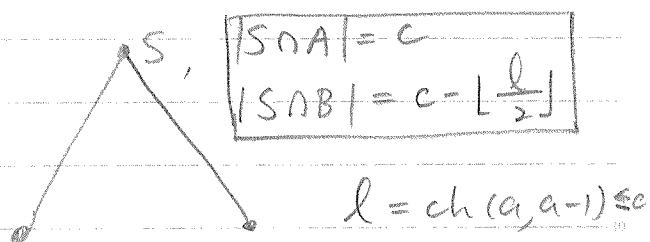
Proof.

Case 1. $b = a-1 \geq 1$

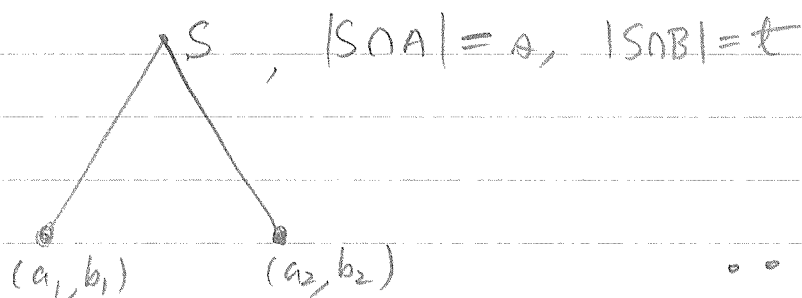
a is even, $a = 2c$



a is odd, $a = 2c+1$



Case 2 $b > a - 1 \geq 1$



The idea of handling more lies, say "three".

Then, each state can be represented by a 4-tuple (a, b, c, d)

where $|A| = a$, $|B| = b$, $|C| = c$ and $|D| = d$.

A: $e \in A$ iff none of the answers is a lie (Truth set)

B: $e \in B$ iff exactly one of the answers is a lie (One lie set)

C: $e \in C$ iff exactly two of the answers are lies (Two lies set)

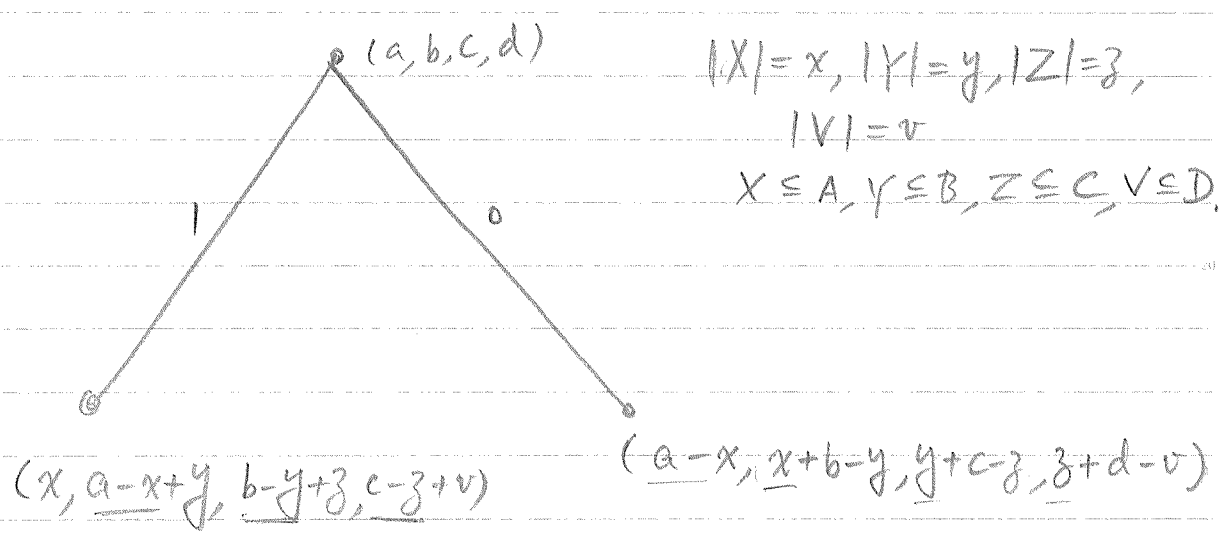
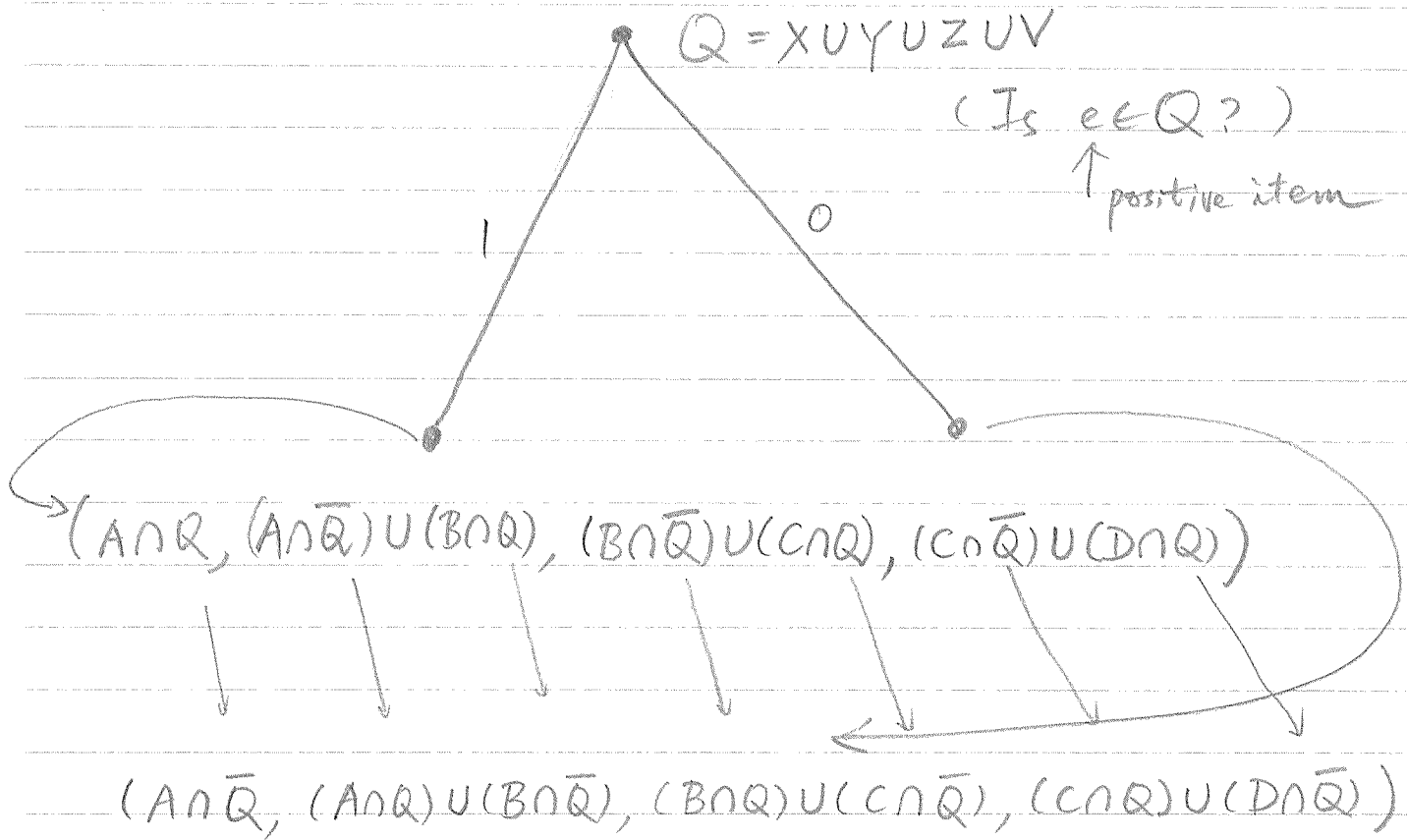
D: $e \in D$ iff exactly three of the answers are lies (Three lies set)

✓ positive item (real)

Queries: "Is $e \in Q$?" where $Q = XUYUZUV \subseteq N$,

and $X \subseteq A$, $Y \subseteq B$, $Z \subseteq C$ and $V \subseteq D$.

(A, B, C, D) ← state of sets



Weight of (a, b, c, d) : with k more tests to do.

$$(a \cdot \binom{k}{3} + b \cdot \binom{k}{2} + c \binom{k}{1} + d) = w_k(a, b, c, d)$$

In truth set, there are three lies to come!

In lie sets, there are 2, 1, 0 lies to come!

Definition

The character of the state (a, b, c, d) ,

$$ch(a, b, c, d) = \min \{ k \mid w_k(a, b, c, d) \leq 2^k \}$$

Similar but more complicated arguments are needed in obtaining the answer.

So, find $M^3(1, 10^6)$.

Please refer to :

A. Negro and Sereno,

1. Solution of Ulam's problem on binary search with three lies,

JCT(A) 59 (1992), 149-150.

2. An Ulam's searching game with three lies, Advances

in Applied Math., Vol. 13, 4, 1992, 404-428.

Two lies

W. Guzicki, Ulam's searching game with two lies,

JCT(A) 54, 1990.

Now, we come back to consider non-adaptive algorithms.

Lies are corresponding with errors occurred in outcome vector.

If we have an error, then a $(1;3)$ -disjunct matrix can

be applied to find a positive with one error. So, if the Ulam's

game is concerned, it is expected to find a $t \times 10^6$ matrix

such that for any two columns C_i and C_j , $|C_i \setminus C_j| \geq 3$.

In case that we can find a positive in 25 tests, then

it is expected that we can let $X = \mathbb{Z}_{25}$ and there are

10^6 subsets of \mathbb{Z}_{25} , such that for any two of them C_i and C_j ,

$|C_i \setminus C_j| \geq 3$.

Exercise 1. Can the above question be done? If not how many

elements in X we need?

Exercise 2. In order to find a positive with 3 errors, then we

need a $(1;7)$ -disjunct matrix. Can we estimate how many tests

we have in order to solve the problem $M^3(1, 10^6)$?

The construction of $(1;3)$ -disjunct matrices

Observation

1. Let (X, \mathcal{B}) be a block design: $2-(v,4,1)$ design. Then we have a $(1;3)$ -disjunct matrix: $v \times b$ $(0,1)$ -matrix where $b = v(v-1)/2$.
2. If $v=25$, then a $2-(25,4,1)$ design exists. Hence, we have a 25×50 $(1;3)$ -disjunct matrix.
3. If the block size we use in \mathcal{B} is k , then two blocks in \mathcal{B} can have at most $k-3$ elements in common.
So, if $X = \mathbb{Z}_{25}$, how many blocks we can get which satisfies property 3?