

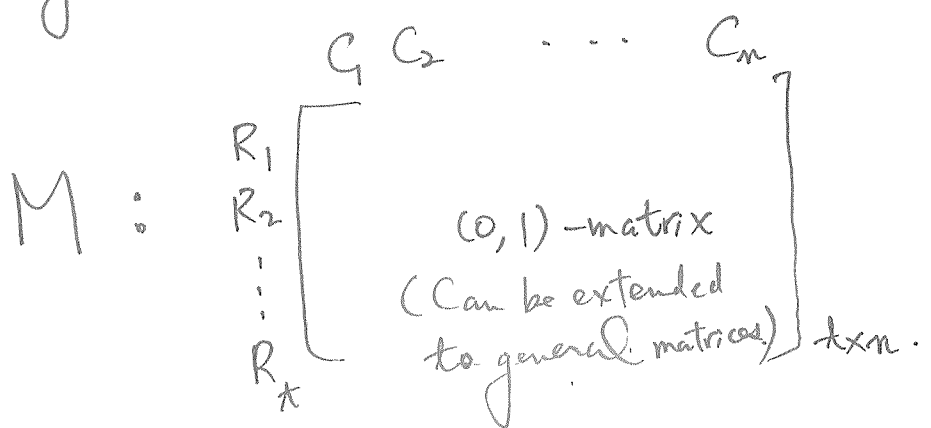
Definition (Non-adaptive)

A group testing algorithm is non-adaptive if all tests must be specified without knowing the outcomes of other tests.

(Fact)

A nonadaptive algorithm can be represented by a $t \times n$ $(0,1)$ -matrix if t tests are specified for a group testing on n items. (Matrix-representation)

Let M be the $(0,1)$ -matrix, R_i and C_j denote row i and column j . We can view R_i as a set of column indices corresponding to the 1-entries.



Definitions

(2)

1. Let A and B are two sets in $[1, n]$. Then the Boolean sum of A and B , denoted $A \oplus B$, is the set of elements in A or B .

2. We shall call $A \oplus B$ as the union of A and B .

3. In I_n , each C_j is the set of row indices i such that $m_{ij} = 1$ in M .

($m_{ij} = 1$)

4. Each C_j can be viewed as a set (subset of $[1, n]$).

5. M is called d -separable (\bar{d} -separable) if the unions of d columns (up to d columns) are all distinct.

6. M will be called d -disjunct if the union of any d columns does not contain any other column.

Group Testing and Superimposed Codes

③

(Non-adaptive)

(*) A $t \times n$ d -separable (d -separable or d -disjunct) matrix generates a nonadaptive (d, n) algorithm of group testing with t tests by associating the columns with items, the rows with tests, and interpreting a 1-entry in cell (i, j) ($m_{ij} = 1$) as the containment (item j is in pool i), and the 0-entry a noncontainment of item j in test (pool) i .

(**) Consider the set of columns C_1, C_2, \dots, C_n as n subsets of $[1, t]$. (is a superimposed code)
 $(M = \{C_1, C_2, \dots, C_n\})$ with length t , volume n , distance D and strength d if $\forall S \subseteq [1, n]$, with $|S| \leq d$ and $i \in [1, n] \setminus S$,

$$|C_i - \bigcup_{j \in S} C_j| \geq D.$$

④

$D=1, \Rightarrow d$ -disjunct

Let $z = D-1$.

A matrix ($t \times n$ incidence matrix) of a superimposed code M with distance D and strength d is called a (d, z) -disjunct matrix.

$(d, 0)$ -disjunct $\approx d$ -disjunct.

How to decode? (nonadaptive algorithm)

Requirements:

(1) d -separable

The unions of any d columns (in $\binom{[1, n]}{d}$) must have distinct outcomes.

Let S, S' be two distinct d -subset of $[1, n]$.

② $\bigcup_{x \in S} C_x \neq \bigcup_{y \in S'} C_y \Rightarrow$ The group testing has different outcomes!

(?)

(5)

(2) \bar{d} -separable

For $|S| \leq d, |S'| \leq d,$

$$\bigcup_{x \in S} C_x \neq \bigcup_{y \in S'} C_y.$$

For example, if $n=20, d=3$, then we need

$\binom{20}{3} + \binom{20}{2} + \binom{20}{1} + \binom{20}{0}$ distinct outcome vectors.

This fact provides a lower bound of the number of tests. Let t be the # of tests. Then

$$2^t \geq \binom{20}{3} + \binom{20}{2} + \binom{20}{1} + \binom{20}{0}.$$

$$\begin{array}{r} \downarrow \\ 3 \times 19 \times 20 \quad 19 \times 10 \quad 20 \quad 1 \\ 1140 + 190 + 20 + 1 \\ = 1351. \end{array}$$

$$t \geq 11 !$$

(3) d -disjunct $\Rightarrow \bar{d}$ -separable

(6)

Proof.

two distinct subset of $[1, t]$

Suppose not. There exist S and S' such that
 $1 \leq |S| \leq d, |S'| \leq d$, and $\bigcup_{x \in S} C_x = \bigcup_{y \in S'} C_y$. Let $i \in S \setminus S'$

Then $C_i \subseteq \bigcup_{x \in S} C_x = \bigcup_{y \in S'} C_y$. Therefore, M is not d -disjunct.

Disjunctness \Rightarrow Easy decoding

$$\left[\begin{array}{c} \\ \\ \\ \\ \\ \\ \\ \end{array} \right] \begin{pmatrix} z_1 \\ z_2 \\ \vdots \\ z_n \end{pmatrix}$$

↑
outcome vector

If $z_i = \begin{cases} 0, & \text{then } R_i \text{ is called a negative pool,} \\ 1, & \text{then } R_i \text{ ————— positive pool.} \end{cases}$

(***)

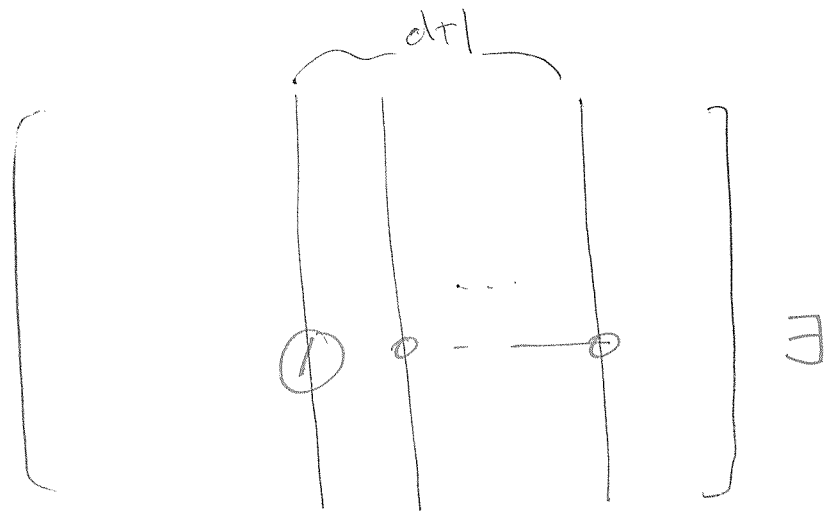
(Fact)


⑦

If $z_i = 0$, then $\forall x \in R_i$, the item corresponds to C_x is negative.

(Fact)

If M is d -disjunct, then the union of zero pools corresponds to at least $n-d$ pure (good) items.



If this item is negative, then it looks like . Not possible!