

The idea of maximal counterexample

10, 2

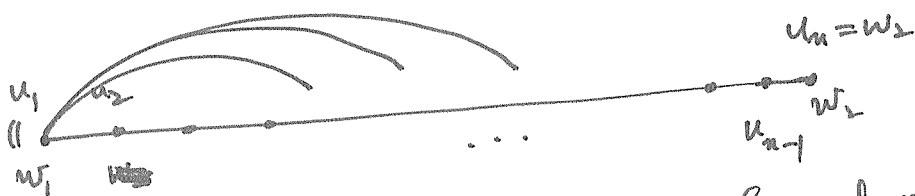
Theorem (Ore, 1960 AMM)

If a graph of order $n \geq 3$ such that for all distinct nonadjacent vertices u and v , $\deg(u) + \deg(v) \geq n$, then G is hamiltonian.

Proof: Assume the theorem is false. Hence there exists a "maximal nonhamiltonian" graph G of order $n \geq 3$ that satisfies the assumption of the theorem. This implies that

$\forall w_1 \neq w_2$, $G + w_1 w_2$ is hamiltonian. (1)

(Note that such pair of vertices w_1 and w_2 does exist since G is not hamiltonian.)



By (1), we have a hamiltonian path with end vertices w_1 and w_2 , see above figure. Now, consider the neighbors of

w_1 and w_2 respectively. Let $N(w_1) = \{u_{i_1}, u_{i_2}, \dots, u_{i_k}\}$
with $i_1 < i_2 < \dots < i_k \leq n-1$. Clearly, $\{u_{i_1-1}, u_{i_2-1}, \dots, u_{i_k-1}\}$

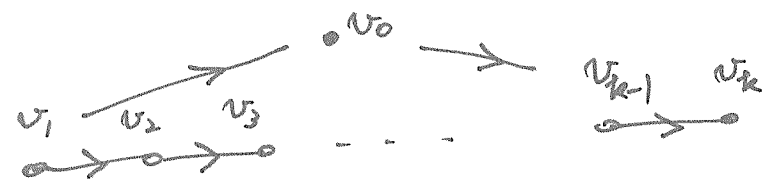
$\cap N(w_2)$ is an empty set, so, $|N(w_1)| + |N(w_2)| \leq n-1$. (?)
 $\deg(w_1) \quad \deg(w_2) \quad \rightarrow \leftarrow$

Theorem (Tournament's fact) (Rédei, 1934)

Every tournament contains a hamiltonian path.
(digraph)

Proof:

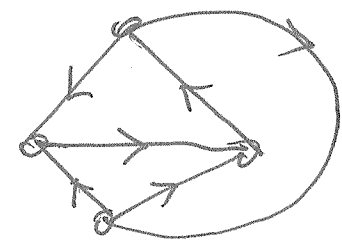
Let $\langle v_1, v_2, \dots, v_k \rangle$ be the longest path in a tournament T . See the following figure for reference.



If $k = |T|$, then we are done.

Otherwise, assume that $k < |T|$. Then, $\exists v_0 \in V(T)$ which is not on $\langle v_1, \dots, v_k \rangle$. By assumption, (v_1, v_0) and (v_0, v_k) are arcs in T . This also implies that $\exists 2 \leq i \leq k-1$, such that (v_i, v_0) and (v_0, v_{i+1}) are arcs in T . In this case, we have a longer path, namely, $\langle v_1, v_2, \dots, v_i, v_0, v_{i+1}, \dots, v_k \rangle$.

Say some more about tournaments! Can you?



Connectivity

Definition (Vertex-connectivity)

The vertex-connectivity or simply connectivity $\kappa(G)$ of a graph G is the minimum number of vertices whose removal from G results in a disconnected or trivial graph.

Definition (Edge-connectivity)

The edge-connectivity $\kappa_1(G)$ of a graph G is the minimum number of edges whose removal from G results in a disconnected (or trivial) graph.

Theorem \forall graph G , $\kappa(G) \leq \kappa_1(G) \leq \delta(G)$.

Proof. (Easy to see?)

Question $\forall 1 \leq a \leq b \leq c$; construct a graph G such that $\kappa(G) = a$, $\kappa_1(G) = b$ and $\delta(G) = c$.