

1. Prove that $ex(n; K_{r+1}) = \|T_r(n)\|$ and $T_r(n)$ is the unique extremal graph.

Proof: Since $T_r(n)$ does not contain a subgraph K_{r+1} , $ex(n; K_{r+1}) \geq \|T_r(n)\|$.

Claim: $K_{r+1}, ex(n; K_{r+1}) \leq \|T_r(n)\|$.

Method 1:

Lemma: Let G be an r -partite graph of order n . then $\|G\| \leq \|T_r(n)\|$.

By lemma, it suffices to show that if G does not contain a subgraph K_{r+1} , then there exist an r -partite graph H s.t. $\|G\| \leq \|H\|$.

By **induction on r** .

It is clear that $r=1$ is true.

Let G be a graph with $\Delta(G)$ and G does not a subgraph K_{r+1} .

Let $x \in V(G)$ s.t. $\deg_G(x) = \Delta(G)$ and $G' = \langle N_G(x) \rangle_G$.

Clearly, G' does not contain K_r as a subgraph (如果有的話, G 就包含 K_{r+1}).

By induction hypothesis, there exists an $(r-1)$ -partite H' s.t. $\|G'\| \leq \|H'\|$.

Let $S = V(G) \setminus V(H')$. Then the graph $H' \vee S$ is an r -partite graph with $\|H'\| + [n - \Delta(G)] \Delta(G)$ edges.

Hence $\|T_r(n)\| = \|H'\| + [n - \Delta(G)] \Delta(G) > \|G'\| + \sum_{v \in S} \deg_G(v) \geq \|G\|$.

Proof 2: Let G be the extremal graph of order n . $\Rightarrow \|G\| = ex(n, K_{r+1})$.

Induction on n .

<i> $n \leq r$. It is clear that G is K_r and $T_r(n) = \|K_r\| = \|G\|$.

<ii> $n > r$. Claim: G does not contain $K_{r+1} \Rightarrow G$ contain K_r .

Proof of claim: If making x and y adjacent to create K_{r+1} , the other $r-1$ vertices must

be all adjacent to x and y . Thus x or y forms a K_r with the $r-1$ vertices.

$\therefore G$ does not contain $K_{r+1} \Rightarrow G$ contain K_r .

Let $V(K_r) = W$ and $V(G) \setminus W = U$, G' be the graph of U .

$\therefore G$ does not contain K_{r+1} .

\therefore Every vertex not in W has at most $r-1$ neighbor in W .

$$\Rightarrow \|E(G)\| \leq \binom{r}{2} + (r-1)(n-r) + \|\langle U \rangle_G\|$$

$\therefore G'$ does not contain $K_{r+1} \Rightarrow \|\langle U \rangle_G\| \leq \|T_r(n-r)\|$.

$$\text{Then } \|E(G)\| \leq \binom{r}{2} + (r-1)(n-r) + \|T_r(n-r)\| = T_r(n).$$

Proof of uniqueness: Let $\|G\| = \|T_r(n)\|$ and G be an extremal graph.

Let $y \in V(G)$, $\deg_G(y) = \delta(G)$.

Then $\|G-y\| = \|G\| - \delta(G)$. (By induction).

$G-y$ must be isomorphic to $T_r(n-1)$.

\Rightarrow Then the smallest partite set in $G-y$ has $\lfloor \frac{n-1}{r} \rfloor$ vertices.

$$\text{Hence } |\langle y, G-y \rangle| = (n-1) - \lfloor \frac{n-1}{r} \rfloor = n - \lfloor \frac{n}{r} \rfloor.$$

$\therefore G-y$ is $T_r(n-1)$.

$\therefore y$ 最少跟一個 partite set is non-adjacent. (否則會有 K_{r+1})

And the partite set must be of size $\lfloor \frac{n-1}{r} \rfloor$. Hence $G \cong T_r(n)$.

2. Prove that $ex(n; K_{s,t}) \leq \frac{1}{2}(s-1)^{\frac{1}{t}} n^{2-\frac{1}{t}} + \frac{1}{2}(t-1)n$, $t \leq s$.

(Give at least two different proofs.)

pf:

(First proof)

Let G be a graph of order n without $K_{s,t}$ with a maximal number of edges.

Let the degree sequence of G be $\langle d_1, d_2, \dots, d_n \rangle$, $d_1 \leq d_2 \leq \dots \leq d_n$.

Consider the t -stars in G .

Since G does not contain $K_{s,t}$, there are at most $\binom{n}{t}^{(s-1)}$ t -stars in G .

Let the average degree of G be $d = \frac{2\|G\|}{n}$.

Then we have $\binom{n}{t}^{(s-1)} \geq \sum_{i=1}^n \binom{d_i}{t} \geq n \binom{d}{t}$.

$$\Rightarrow (s-1)n^{-1} \geq \frac{\binom{d}{t}}{\binom{n}{t}} = \frac{d(d-1)\dots(d-t+1)}{n(n-1)\dots(n-t+1)} \geq \left(\frac{d-t+1}{n}\right)^t.$$

$$\Rightarrow (s-1)^{\frac{1}{t}} n^{-\frac{1}{t}} \geq \frac{d-t+1}{n}.$$

$$\Rightarrow \frac{2\|G\|}{n} = d \leq (s-1)^{\frac{1}{t}} n^{1-\frac{1}{t}} + (t-1).$$

$$\Rightarrow \|G\| \leq \frac{1}{2}(s-1)^{\frac{1}{t}} n^{2-\frac{1}{t}} + \frac{1}{2}(t-1)n.$$

(Second proof)

Lemma :

Let $2 \leq s \leq m$, $2 \leq t \leq n$, $0 \leq r < m$, $z = km+r$ and $z = my$.

Let $G_2(m,n)$ be a graph of size z that does not contain $K_{s,t}$.

$$\text{Then } m \binom{y}{t} \leq (m-r) \binom{k}{t} + r \binom{k+1}{t} \leq (s-1) \binom{n}{t}.$$

pf:

Let $G_2(m,n) = (V_1, V_2)$, where $|V_1| = m$ and $|V_2| = n$.

Define $H = \left(V_1, \binom{V_2}{t} \right)$, where $\binom{V_2}{t}$ is the collection of all t -subsets of V_2 .

$$\forall x \in V_1, A \in \binom{V_2}{t}, x \square A \text{ iff } x \square a, \forall a \in A.$$

$$\Rightarrow \|H\| = \sum_{x \in V_1} \binom{\deg_G(x)}{t}.$$

Note that $\forall B \in \binom{V_2}{t}, \deg_G(B) \leq s-1$. (Otherwise $K_{s,t} \leq G$.)

$$\Rightarrow \|H\| = \sum_{x \in V_1} \binom{\deg_G(x)}{t} \leq (s-1) \binom{n}{t}.$$

Since $\sum_{x \in V_1} \deg_G(x) = \|G\| = z = km+r = my$, we have

$$m \binom{y}{t} \leq (m-r) \binom{k}{t} + r \binom{k+1}{t} \leq \sum_{x \in V_1} \binom{\deg_G(x)}{t} \leq (s-1) \binom{n}{t}.$$

Theorem :

$$z(m, n; s, t) \leq (s-1)^{\frac{1}{t}} (n-t+1) m^{1-\frac{1}{t}} + (t-1)m.$$

pf:

$$\text{By the above lemma, we have } m \binom{y}{t} \leq (s-1) \binom{n}{t}.$$

$$\Rightarrow \frac{\binom{y}{t}}{\binom{n}{t}} \leq (s-1) m^{-1}.$$

$$\Rightarrow \frac{y(y-1)\dots(y-t+1)}{n(n-1)\dots(n-t+1)} \leq (s-1) m^{-1}.$$

Since $G \leq K_{m,n}$, $my = z = \|G\| \leq \|K_{m,n}\| = mn$ and $y \leq n$.

$$\Rightarrow \frac{(y-t+1)(y-t+1)\dots(y-t+1)}{(n-t+1)(n-t+1)\dots(n-t+1)} \leq \frac{y(y-1)\dots(y-t+1)}{n(n-1)\dots(n-t+1)} \leq (s-1) m^{-1}.$$

$$\Rightarrow \frac{(y-t+1)^t}{(n-t+1)^t} \leq (s-1) m^{-1}.$$

$$\Rightarrow (y-t+1)^t \leq (s-1)(n-t+1)^t m^{-1}.$$

$$\Rightarrow y-t+1 \leq (s-1)^{\frac{1}{t}} (n-t+1) m^{-\frac{1}{t}}.$$

$$\Rightarrow \frac{z}{m} = y \leq (s-1)^{\frac{1}{t}} (n-t+1) m^{-\frac{1}{t}} + (t-1).$$

$$\Rightarrow z \leq (s-1)^{\frac{1}{t}} (n-t+1) m^{1-\frac{1}{t}} + (t-1)m.$$

Now we prove this exercise.

By the above theorem, we have

$$z(m, n; s, t) \leq (s-1)^{\frac{1}{t}}(n-t+1)m^{1-\frac{1}{t}} + (t-1)m.$$

$$\Rightarrow z(n, n; s, t) \leq (s-1)^{\frac{1}{t}}(n-t+1)n^{1-\frac{1}{t}} + (t-1)n \leq (s-1)^{\frac{1}{t}}n^{2-\frac{1}{t}} + (t-1)n.$$

Hence it suffices to claim $ex(n; K_{s,t}) \leq \frac{1}{2}z(n, n; s, t)$.

Let G be a graph without $K_{s,t}$ with a maximal number of edges and

$$V(G) = \{x_1, x_2, \dots, x_n\}.$$

Define $H = (V_1, V_2)$ such that $V_1 = \{x'_1, x'_2, \dots, x'_n\}$, $V_2 = \{x''_1, x''_2, \dots, x''_n\}$ and

$$x'_i \square x''_j \text{ iff } x_i \square x_j.$$

Since G does not contain $K_{s,t}$, we have H does not contain $K_{s,t}$, and hence

$$\|H\| \leq z(n, n; s, t).$$

By the construction of H , we have $\|H\| = 2\|G\|$.

$$\Rightarrow ex(n; K_{s,t}) = \|G\| \leq \frac{1}{2}\|H\| \leq \frac{1}{2}z(n, n; s, t) \leq \frac{1}{2}(s-1)^{\frac{1}{t}}n^{2-\frac{1}{t}} + \frac{1}{2}(t-1)n. \quad \square$$

3. Find a graph G of order 100 s.t. G does not contain $K_{2,2}$ and $\|G\| \approx 500$

Proof. Since $ex(n; K_{2,2}) \leq \frac{n}{4}(1 + \sqrt{4n-3})$ and equality holds whenever $n = q^2 + q + 1$, q is a prime power. Let $q = 3^2$, then

$$n = 91 \quad \text{and} \quad ex(91; K_{2,2}) = \frac{91}{4}(1 + \sqrt{4 \cdot 91 - 3}) = 455$$

Therefore exist a graph G' does not contain $K_{2,2}$ with $\|G'\| = 455$ and $|G'| = 91$. Let $G = G' \cup C_9$. Since G' and C_9 does not contain $K_{2,2}$, therefore

$$G \text{ does not contain } K_{2,2} \quad \text{and} \quad |G| = 100, \|G\| = 455 + 9 = 464$$

□

4. Find a C_4 -saturated graph of order n with minimum number of edges for each positive integer n .

First, we construct C_4 -saturated graph with $\left\lfloor \frac{3n-5}{2} \right\rfloor$ edges.

Second, we prove such graph with minimum number of edges for each integer n .

Now, we construct C_4 -saturated graph with order n .

Step 1:

Take $t = \left\lfloor \frac{n-3}{2} \right\rfloor$ triangle sharing a vertex w .

Step2.

$$\text{Join } p = n - 2t - 1 = n - 2 \left\lfloor \frac{n-3}{2} \right\rfloor - 1 = \begin{cases} 3, n = 2k + 4 & k \in \{0\} \cup \mathbb{Z}^+ \\ 2, n = 2k + 3 & k \in \{0\} \cup \mathbb{Z}^+ \end{cases}$$

new vertices to p vertices of one of these triangles using independent edges, (i.e. the new vertices have no common neighbors.)

So, we want to prove the graph we obtained is C_4 -saturated.

(case1)

(a) If add $\{p_1, p_2\}$ to be new edge:

Then, clearly $p_1 - a_1 - a_2 - p_2 - p_1$ is C_4 , where a_1, a_2 are vertices of one of the t triangles.

(b) Otherwise, add any edge $\{x, z\}$ $\{x, z\} \neq \{p_1, p_2\}$:

Then, for $x \in V(G)$, $|N(x) \cap N(y)| = 1$.

Moreover, $\{z, w\} \in E(G)$.

So, let $y \in N(x) \cap N(w)$

Then $x-z-w-y-x$ is C_4 .

Hence, by (a) (b), when $n=2k+3$, graph is C_4 -saturated.

(case2)

$$n=2k+4, k \in \{0\} \cup \mathbb{Z}^+.$$

Let the new three vertices p_1, p_2, p_3 , where $|N(p_3) \cap w| = 0$.

Consider two cases:

(a) If add $\{x, z\}$ to be new edge, $x, z \neq p_3$.

Then, we can use the same way of case1 to discuss.

(b) If x or z is p_3 , then w.l.o.g let $x = p_3$.

Then, except p_3 ,

$$\forall x_1 \in V(G), |N(x_1) \cap N(w)| = 1$$

Thus, $|N(z) \cap N(w)| = 1$.

Let $y \in N(z) \cap N(w)$

Then $y-z-x-w-y$ is C_4 .

Hence, by (a) (b), when $n=2k+4$, graph G is C_4 - saturated.

Hence, by case1 and case2, the graph we constructed is C_4 - saturated.

Moreover, the edges of graph are $\left\lfloor \frac{3n-5}{2} \right\rfloor$.

Then, after we constructing C_4 - saturated graph, let's

prove its size is minimum.

Let P_k be a path on k edges.

A graph is P_k -connected if all pairs of non-adjacent vertices are P_k -connected.

So, if adding an edge between two vertices creates a C_k , then a P_{k-1} must connect the two vertices.

Thus, a C_k -saturated graph is P_{k-1} -connected.

Now, we want to prove the following results:

Let G be a P_3 -connected graph, then $\|G\| \geq \frac{3(n-2)}{2}$.

(Lemma)

Let G be a connected graph. If each edge of G is in C_3 , then $\|G\| \geq \frac{3(n-1)}{2}$.

(pf)

Since G is connected, G has a spanning tree T , $\|T\| \geq n-1$.

An edge of $G-T$ can form a C_3 with at most two edges of T .

So, $\|G-T\| \geq \frac{1}{2}(n-1)$.

Thus, $\|T\| \geq (n-1) + \frac{1}{2}(n-1) = \frac{3}{2}(n-1)$.

(Thm)

Let G be a P_3 -connected graph. Then $\|G\| \geq \frac{3}{2}(n-2)$.

(pf)

It is true for $n=1$.

So, assume $n \geq 2$.

Since G is connected, $\delta > 0$.

If $\delta \geq 3$, then $\|G\| \geq \frac{3}{2}n > \frac{3}{2}(n-2)$.

So, assume $\delta=1$ or $\delta=2$.

(case1)

$$\delta=1$$

Let $\deg(v)=1$ and $N(v)=\{w\}$.

Let T be the breath-first tree of G rooted at v .

$$\text{Let } X=\{x|d(x,v)=2\} \quad Y=\{x|d(x,v)=3\}$$

Since G is P_3 -connected, all vertices are distance at most 3 from v .

So, all vertices other than v and w are in X or Y .

$$\text{Thus, } |X|+|Y|=n-2.$$

Because $x \in X$ is not adjacent to v , there is a path $v-w-q-x$.

Since $d(q,x)=2$, $q \in X$.

Hence, $q \in X$, $N(q) \cap X \neq \emptyset$.

$$\text{Let } Y_1 = \{z|\deg(z)=1, z \in V(Y)\}, |Y_1|=k$$

$$\text{Let } Y_2 = Y - Y_1.$$

Let X_1 be the vertices of X adjacent to Y_1 , and let $X_2=X-X_1$

Then, for all $y_1, y_2 \in Y_1$, since G is P_3 -connected, there is a path $y_1 - x_1 - x_2 - y_2$, where $x_1, x_2 \in X_1$.

Moreover, each vertex in X_1 , is adjacent to exactly one vertex in Y_1 , and every pair of vertices in X_1 are adjacent.

Thus, $|X_1|=|Y_1|=k$, and X_1 is clique.

Since each vertices in X_2 and Y_2 is incident to at least one edge not in T , the number of edges noy in T

$$\text{is } \|G-T\| \geq \binom{|X_1|}{2} + \frac{|X_1|+|X_2|}{2} = \binom{k}{2} + \frac{n-2-2k}{2} \geq \frac{n-4}{2}.$$

$$\rightarrow \|G\| \geq \frac{1}{2}(n-4) + n-1 = \frac{3}{2}(n-2).$$

(case2)

$\delta=2$ and $\exists v$ is not in C_3 , where $\deg(v)=2$.

Let T be the breath-first tree of G rooted at v .

$$\text{Let } X=\{x|d(x,v)=2\} \quad Y=\{x|d(x,v)=3\}$$

Since a P_3 connects v to x , $x \in X_1$

There is an edge incident to x which is not in T .

Further, since $\delta=2$, each $y \in Y$ is also incident to an edge not in T .

$$\text{So, } \|G-T\| \geq \frac{1}{2}(|X|+|Y|) = \frac{1}{2}(n-3).$$

$$\text{Therefore, } \|G\| \geq \frac{1}{2}(3n-5).$$

(case3)

$\delta=2$ and $\forall v$ is in C_3 , where $\deg(v)=2$.

Let S be the subgraph of all C_3 's with one or more degree 2 vertices.

Let S_i , $1 \leq i \leq m$ be the components of S .

By Lemma 1, $\|S_i\| \geq \frac{3}{2}(|S_i| - 1) \forall i$.

Moreover, since the degree 2 vertices of two components are P_3 -connected, there is an edge between every pair of components of S .

So, there are at least $\binom{m}{2}$ edges between components of S .

Then, the number of edges in the induced subgraph induced by the vertices of S is $\| \langle S \rangle \| \geq \sum_{i=1}^m \|S_i\| + \binom{m}{2} \geq$

$$\sum_{i=1}^m \frac{3}{2}(|S_i| - 1) + \binom{m}{2} = \frac{3}{2}|S| + \frac{m^2}{2} - 2m \geq \frac{3}{2}|S| - 2.$$

Since all degree 2 vertices are in S , vertices not in S have 3 or more.

$$\text{Thus, } \|G\| \geq \frac{3}{2}|S| - 2 + \frac{3}{2}(|G| - |S|) = \frac{3}{2}n - 2.$$

Hence, by case1, case2, case3, we can conclude $\|G\| \geq \frac{3}{2}(n - 2)$.

$$\text{Thus, } \|G\| \geq \left\lceil \frac{3n-6}{2} \right\rceil = \left\lfloor \frac{3n-5}{2} \right\rfloor.$$

So, the graph we construct is the graph of minimum size.

Graph Theory

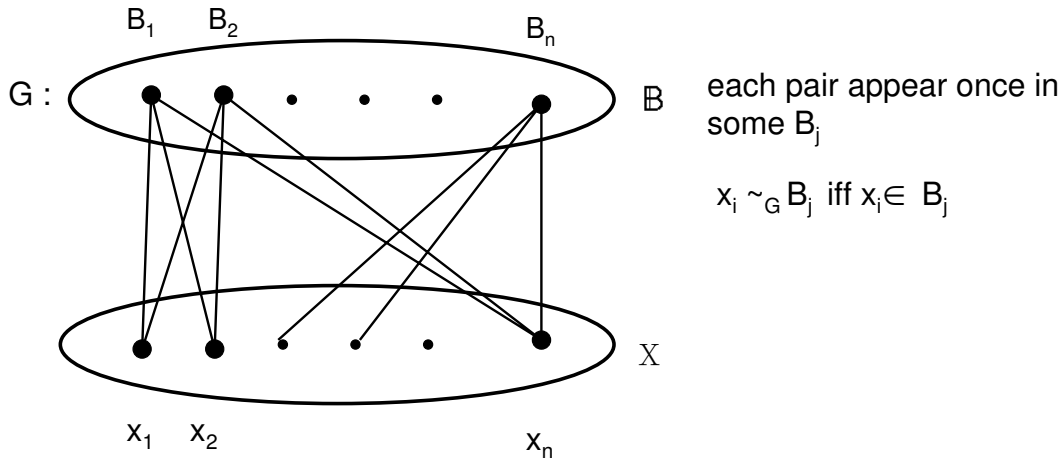
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6. Find $z(n, n; 2, 2)$ for as many n as possible.

sol:

When $n = q^2 + q + 1$ which q is prime power, $z(n, n; 2, 2) \leq \frac{n}{2}[1 + (4n - 3)^{\frac{1}{2}}]$ holds.

Since q is prime power, there exists a projective plane, let $G = (\mathbf{X}, \mathbb{B})$ be $||G|| = (q + 1)(q^2 + q + 1)$. And since G is projective plane, it doesn't contain $K_{2,2}$. Since $C_4 \cong K_{2,2}$, and G is $2-(q^2 + q + 1, q + 1, 1)$ design.



Since $n = q^2 + q + 1$, it implies $q^2 + q + (1 - n) = 0$, we obtain $q = \frac{-1 \pm \sqrt{1 - 4(1 - n)}}{2} = \frac{-1 \pm (4n - 3)^{\frac{1}{2}}}{2}$ (we choose positive). Thus, $||G|| = (q + 1)(q^2 + q + 1) = \frac{1 + (4n - 3)^{\frac{1}{2}}}{2} \cdot n$. We obtain $z(n, n; s, t) \leq \frac{n}{2}[1 + (4n - 3)^{\frac{1}{2}}]$ as required.

7. Find a subgraph G of $K_{9,12}$ which is C_4 -saturated and $\|G\| = 28$ (Explain your answer).

Definition: (1). A partial $2 - (v, K, 1)$ design (\mathbb{X}, \mathbb{B}) is said to be an extension of a $2 - (v, K, 1)$ design $(\mathbb{X}, \mathbb{B}')$ if $\mathbb{B} \neq \mathbb{B}'$, $|\mathbb{B}| = |\mathbb{B}'|$ and $B' \subseteq B$ for each $B \in \mathbb{B}$ and $B' \in \mathbb{B}'$.

(2). A partial $2 - (v, K, 1)$ design (\mathbb{X}, \mathbb{B}) is said to be non-extendable if it has no extension.

Lemma 1 The variety-block graph $G_{(\mathbb{X}, \mathbb{B})}$ of a partial $2 - (v, K, 1)$ design $2 - (v, K, 1)$ is a C_4 -saturated graph of $K_{|\mathbb{X}|, |\mathbb{B}|} \Leftrightarrow (\mathbb{X}, \mathbb{B})$ is non-extendable.

Proof of Lemma 1:

(\Rightarrow) Let the variety-block graph $G_{(\mathbb{X}, \mathbb{B})}$ of a partial $2 - (v, K, 1)$ design (\mathbb{X}, \mathbb{B}) be a C_4 -saturated graph of $K_{|\mathbb{X}|, |\mathbb{B}|}$. If (\mathbb{X}, \mathbb{B}) is extendable then let $(\mathbb{X}, \mathbb{B}')$ be the extension of (\mathbb{X}, \mathbb{B}) and $x \in B' \setminus B$ where $B' \subseteq B$, $B \in \mathbb{B}$ and $B' \in \mathbb{B}'$. Since $(\mathbb{X}, \mathbb{B}')$ is also a partial $2 - (v, K, 1)$ design, adding the edge xB to $G_{(\mathbb{X}, \mathbb{B})}$ won't make C_4 . Hence $G_{(\mathbb{X}, \mathbb{B})}$ is not a C_4 -saturated ($\rightarrow \leftarrow$). Thus (\mathbb{X}, \mathbb{B}) is non-extendable.

(\Leftarrow) Suppose (\mathbb{X}, \mathbb{B}) is non-extendable. If $G_{(\mathbb{X}, \mathbb{B})}$ is not C_4 -saturated, then there exists an edge such that adding the edge to $G_{(\mathbb{X}, \mathbb{B})}$ is also C_4 -free. W.L.O.G. we let this edge be xB . Then let \mathbb{B}' be the set of blocks obtained from \mathbb{B} by replacing B with $B' = B \cup \{x\}$. Hence $(\mathbb{X}, \mathbb{B}')$ is an extension of (\mathbb{X}, \mathbb{B}) ($\rightarrow \leftarrow$).

(Since adding xB to $G_{(\mathbb{X}, \mathbb{B})}$ is also C_4 -free, $(\mathbb{X}, \mathbb{B}')$ is also a partial $2 - (v, K, 1)$ design. Moreover, $\mathbb{B} \neq \mathbb{B}'$, $|\mathbb{B}| = |\mathbb{B}'|$ and $B' \subseteq B$ for each $B \in \mathbb{B}$ and $B' \in \mathbb{B}'$. Hence $(\mathbb{X}, \mathbb{B}')$ is an extension of (\mathbb{X}, \mathbb{B})). \square

sol: Let $\mathbb{X} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$$\mathbb{B} = \{\{1, 2, 3, 4, 5, 6, 7\}, \{1, 8, 9\}, \{2, 9\}, \{3, 9\}, \{4, 9\}, \{5, 9\}, \{6, 9\}, \{7, 9\}, \{2, 8\}, \{3, 8\}, \{1\}, \{1\}\} \Rightarrow |\mathbb{B}| = 12$$

It's easy to see that (\mathbb{X}, \mathbb{B}) is a partial $2 - (9, K, 1)$ design, where

$K = \{1, 2, 3, 7\}$. (Since except $(4, 8), (5, 8), (6, 8), (7, 8)$ these four pairs, each pair of \mathbb{X} occurs in exactly one block).

Claim: (\mathbb{X}, \mathbb{B}) is non-extendable.

It suffices to prove that if we add x ($x \in \mathbb{X}$) to any block of \mathbb{B} , then there exists a pair which occurs in at least two blocks.

(1). If we add 8 or 9 to $\{1, 2, 3, 4, 5, 6, 7\}$, then $(i, 8)$ or $(i, 9)$ for some $i = 1, 2, \dots, 7$ will occur in at least two blocks.

(2). If we add j to $\{1, 8, 9\}$, $j = 2, 3, \dots, 7$, then $(1, j)$ will occur in two blocks.

(3). If we add j to $\{i, 9\}$, $i = 2, 3, \dots, 7$, $j \in \mathbb{X} \setminus \{i, 9\}$,
 $\left\{ \begin{array}{l} \text{if } j = 8, \text{ then } (8, 9) \text{ occurs in two blocks;} \\ \text{if } j \neq 8, \text{ then } (i, j) \text{ occurs in two blocks.} \end{array} \right.$

(4). If we add j to $\{i, 8\}$, $i = 2, 3$, $j \in \mathbb{X} \setminus \{i, 8\}$, then (i, j) occurs in at least two blocks.

(5). If we add j to $\{1\}$, $j = 2, 3, \dots, 9$, then $(1, j)$ occurs in two blocks.

\Rightarrow by (1) (5) we have (\mathbb{X}, \mathbb{B}) is non-extendable.

Hence by Lemma 1 we have, $G_{(\mathbb{X}, \mathbb{B})}$ is a C_4 -saturated graph of $K_{9,12}$ and

$$\|G_{(\mathbb{X}, \mathbb{B})}\| = \sum_{B \in \mathbb{B}} |B| = 7 + 3 + 2 \times 8 + 1 \times 2 = 28. \quad \square$$

8. Let $\aleph(G) = r+1$, then H not in $T_r(n)$ (If $H \leq T_r(n)$, then $\aleph(G) \leq r$)
 $\Rightarrow ex(n; H) \geq \|T_r(n)\|$.

Since $H \leq K_{(r+1)(t)}$ for some t , $\|T_r(n)\| \leq ex(n; H) \leq \|T_r(n)\| + \varepsilon n^2$
 (By Erdos and stone's Theorem $ex(n; H) \leq \|T_r(n)\| + \varepsilon n^2$)

$$\begin{aligned} \frac{\|T_r(n)\|}{\binom{n}{2}} &\leq \frac{ex(n; H)}{\binom{n}{2}} \leq \frac{\|T_r(n)\|}{\binom{n}{2}} + \frac{\varepsilon n^2}{\binom{n}{2}} \\ &\leq \frac{\|T_r(n)\|}{\binom{n}{2}} + \frac{2\varepsilon n^2}{n(n-1)} \\ &= \frac{\|T_r(n)\|}{\binom{n}{2}} + \frac{2\varepsilon}{1-\frac{1}{n}} \quad (\|H\| \geq 1 \Rightarrow |H| = n \geq 2 \Rightarrow \frac{1}{1-\frac{1}{n}} \leq 2) \\ &\leq \frac{\|T_r(n)\|}{\binom{n}{2}} + 4\varepsilon \end{aligned}$$

Claim $\lim_{n \rightarrow \infty} \frac{\|T_r(n)\|}{\binom{n}{2}} = 1 - \frac{1}{r}$

$$\begin{aligned} \text{Let } n = qr, \quad \|T_r(n)\| &= \binom{n}{2} - r \binom{q}{2} = \binom{n}{2} - \frac{rq(q-1)}{2} \\ &= \binom{n}{2} - \frac{rq(qr-r)}{2r} = \binom{n}{2} - r \frac{n(n-r)}{2r} \\ &= \binom{n}{2} \left(1 - \frac{1}{r}\right) + \frac{n(r-1)}{2r} \end{aligned}$$

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{\|T_r(G)\|}{\binom{n}{2}} &= \lim_{n \rightarrow \infty} \left(1 - \frac{1}{r}\right) + \frac{(r-1)n}{2r} \times \frac{2}{n(n-1)} = 1 - \frac{1}{r} \\ \Rightarrow \lim_{n \rightarrow \infty} \frac{\|T_r(G)\|}{\binom{n}{2}} &= 1 - \frac{1}{r} = \frac{r-1}{r} = \frac{\aleph(G)-2}{\aleph(G)-1} \end{aligned}$$

9. Given n point in the plane. Prove that there are

most $\frac{1}{\sqrt{2}} n^{\frac{3}{2}} + \frac{n}{4}$ pairs with distance 1.

Proof:把 n 個 points 放在 plane 上

若 2 points x, y 距離為 1.則在 G 中有 xy 邊.

Goal: $|E(G)| \leq \frac{1}{\sqrt{2}} n^{\frac{3}{2}} + \frac{n}{4}$

Claim①: $2\binom{n}{2} \geq \sum_{v \in V(G)} \binom{\deg(v)}{2}$

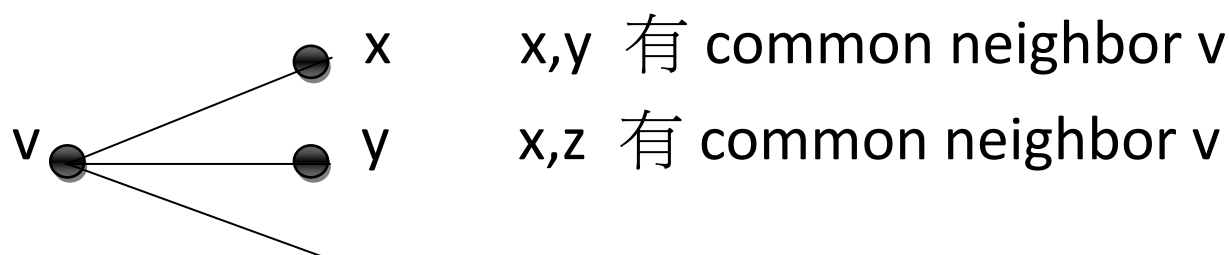
左式為計算 n 點中任兩點的 common neighbor 數.

Since 在平面中任兩點以距離一為半徑作圓, 最多交兩點(相交圓).

意同於,任兩點最多有 2 個 common neighbor 所以最多 $2\binom{n}{2}$ 個.

右式為從圖型上看 common neighbor 數

對每一點的 degree 中任兩點就有 1 common neighbor



● Z

實際上有 $\sum_{v \in V(G)} \binom{\deg(v)}{2}$ 個

$$\Rightarrow 2 \binom{n}{2} \geq \sum_{v \in V(G)} \binom{\deg(v)}{2}$$

Claim ②: $\sum_{v \in V(G)} \binom{\deg(v)}{2} \geq n \binom{\frac{2e}{n}}{2}$

Since $\sum_{v \in V(G)} \deg(v) = 2e$ where $e = |E(G)|$

By $\binom{n}{k}$ is convex function

又右式比左式平均 (i.e degree 的落差可能很大)

$$\Rightarrow \sum_{v \in V(G)} \binom{\deg(v)}{2} \geq n \binom{\frac{2e}{n}}{2}$$

by ① ② $\Rightarrow 2 \binom{n}{2} \geq \sum_{v \in V(G)} \binom{\deg(v)}{2} \geq n \binom{\frac{2e}{n}}{2}$

$$\Rightarrow 2 \frac{n(n-1)}{2} \geq n \frac{\frac{2e}{n}(\frac{2e}{n}-1)}{2}$$

$$\Rightarrow n(n-1) \geq e \left(\frac{2e-n}{n} \right)$$

$$\Rightarrow n^2 - n \geq \frac{2e^2 - en}{n} \Rightarrow 2e^2 - en - (n^3 - n^2) \leq 0$$

$$\Rightarrow e \leq \frac{n + \sqrt{n^2 + 8(n^3 - n^2)}}{2 \cdot 2} = \frac{n}{4} + \frac{\sqrt{8n^3 - 7n^2}}{4}$$

Since

$$\frac{1}{\sqrt{2}} n^{\frac{3}{2}} - \frac{\sqrt{8n^3 - 7n^2}}{4} = \frac{1}{4\sqrt{2}} (4\sqrt{n^3} - \sqrt{16n^3 - 14n^2}) > 0$$

$$\text{Then } \frac{1}{\sqrt{2}} n^{\frac{3}{2}} + \frac{n}{4} \geq \frac{n}{4} + \frac{\sqrt{8n^3 - 7n^2}}{4} \geq e = E(G)$$

\Rightarrow most $\frac{1}{\sqrt{2}} n^{\frac{3}{2}} + \frac{n}{4}$ pairs with distance 1.

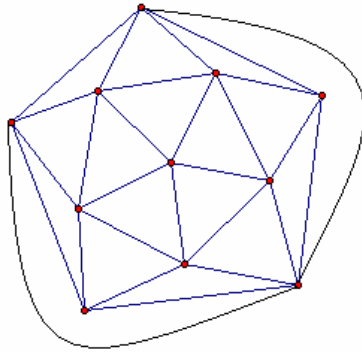
10. Construct a graph G of order 11 such that G does not contain $K_{3,3}$ and G has at least $\lfloor (11^{5/3})/2 \rfloor$ edges.

First of all, we calculate that $\lfloor (11^{5/3})/2 \rfloor = 27$.

(1)

If G is a planar graph, then $|E(G)| \leq 3|V(G)| - 6 = 33 - 6 = 27$.

Thus there exists a planar graph with order 11 and size 27. Moreover, since this graph is planar, it contains no $K_{3,3}$.



This graph is a planar graph since there is no crossing.

It has 11 vertices and 27 edges.

(2)

Let $V(G) = K_1 \cup K_2 \cup K_3 \cup O$, where

$$K_1 = \{K_{11}, K_{12}, K_{13}, K_{14}\}$$

$$K_2 = \{K_{21}\}$$

$$K_3 = \{K_{31}, K_{32}, K_{33}, K_{34}\}$$

$$O = \{O_1, O_2\}, \text{ thus } |V(G)| = 11.$$

Let $\langle K_1 \cup K_2 \rangle \cong \langle K_2 \cup K_3 \rangle \cong K_5$, complete graph of order 5,

$|E_{O, K_2}| = 0$, and O_i is adjacent to O_2 . Moreover, let

$$O_1 \sim K_{11}, K_{12}, K_{31}, K_{32} \text{ and}$$

$$O_2 \sim K_{13}, K_{14}, K_{33}, K_{34}.$$

$$K_{1i} \sim K_{3i}, \forall i = 1, 2, 3, 4.$$

Then the size of graph G would be

$$\|G\| = |K_1 \cup K_2| + |K_2 \cup K_3| + |E_{K_1, K_3}| + |O| + \sum_{i=1}^2 (|E_{K_1, O_i}| + |E_{K_2, O_i}| + |E_{K_3, O_i}|)$$

$$= 10 + 10 + 4 + 1 + 2(2 + 0 + 2) = 33.$$

Claim: G is $K_{3,3}$ -free.

Proof:

First $\langle K_1 \cup K_2 \cup K_3 \rangle$ contains no $K_{3,3}$. If it does contain $G' \cong K_{3,3}$ as a subgraph,

then there must exist $u, v \in V(G')$, $u \in K_1$ and $v \in K_3$ in G .

Then u and v have at most 3 common neighbors, namely K_{21} , the only neighbor u have in K_3 and the only neighbor v have in K_1 . But these 3 vertices have no more common neighbors but u and v , thus forms no $K_{3,3}$.

Then suppose $|V(G') \cap O| \neq \emptyset$.

(i) $|V(G') \cap O| = 1$

Then suppose it's O_1 . There are 4 neighbors of O_1 , besides O_2 .

Any 3 of them have at most one common neighbor, namely K_{21} . Thus forms no $K_{3,3}$.

(ii) $|V(G') \cap O| = 2$

O_1 and O_2 can not be in the same partie in G' , since they only have 2 common neighbors.

Suppose O_1 and O_2 are in different parties in G' . Then $K_{21} \notin V(G')$ since K_2 connects to neither O_1 nor O_2 . In order to create a $K_{3,3}$, there must be 2 vertices that are in the same partie with O_1 in G' which are neighbors of O_1 or O_2 and have at least 2 common neighbors which are also neighbors of O_1 or O_2 .

1) Suppose such 2 vertices, u and v , are in K_1 .

Then they only have one common neighbor fits conditions above.

2) Suppose u is in K_1 , v is in K_3 .

If they are adjacent, then they have no common neighbor fits conditions above.

If they are not adjacent, then they have 2 common neighbors, namely u 's only neighbor in K_3, u' , and v 's only neighbor in K_1, v' , which means if $u = K_{1i}$ for some i , $u' = K_{3i}$, similarly if $v = K_{3j}$, $v' = K_{1j}$ for some j .

However, since O_1, u, v are in same partie in G' , $O_1 \sim u'$ and $O_1 \sim v'$. Thus the nonordered pair (i, j) can only be $(1, 2)$, which means O_2 doesn't connect to either u or v , since $O_2 \sim K_{1x}$ or K_{3x} if and only if $x = 2$ or 3 . Thus this forms no $K_{3,3}$.

And we finish the proof for the claim.