

Graph Theory I Homework III-2 by 陳哲皓

Definition: Let G be a graph and $x \in G$ a vertex, and let G' be obtained from G by adding a vertex x' and joining it to x and all the neighbours of x . We say that G' is obtained from G by expanding the vertex x to an edge xx'

claim1: Any graph obtained from a perfect graph by expanding a vertex is again perfect.

Proof. We use induction on the order of the perfect graph considered. Expanding the vertex of K_1 yields K_2 , which is perfect. For the induction step, let G be a non-trivial perfect graph, and let G' be obtained from G by expanding a vertex $x \in G$ to an edge xx' . For our proof that G' is perfect it suffices to show $\chi(G') \leq \alpha(G')$: every proper induced subgraph H of G' is either isomorphic to an induced subgraph of G or obtained from a proper induced subgraph of G by expanding x ; in either case, H is perfect by assumption and the induction hypothesis, and can hence be coloured with $\omega(H)$ colours.

Let $\omega(G) =: \omega$; then $\omega(G') \in \{\omega, \omega+1\}$. If $\omega(G') = \omega+1$, then

$$\chi(G') \leq \chi(G) + 1 = \omega + 1 = \omega(G')$$

and we are done. So let us assume that $\omega(G') = \omega$. Then x lies in no $K_\omega \subseteq G$: together with x' , this would yield a $K_{\omega+1}$ in G' . Let us colour G with ω colours. Since every $K_\omega \subseteq G$ meets the colour class X of x but not x itself, the graph $H := G - (X \setminus \{x\})$ has clique number $\omega(H) < \omega$. Since G is perfect, we may thus colour H with $\omega - 1$ colours. Now X is independent, so the set $(X \setminus \{x\}) \cup \{x'\} = V(G' - H)$ is also independent. We can therefore extend our $(\omega - 1)$ -colouring of H to an ω -colouring of G' , showing that $\chi(G') \leq \omega = \omega(G')$ as desired \square

claim2: A graph is perfect if and only if its complement is perfect.

Proof. Applying induction on $|G|$, we show that the complement \bar{G} of any perfect graph $G = (V, E)$ is again perfect. For $|G| = 1$ this is trivial, so let $|G| \geq 2$ for the induction step. Let \mathcal{K} denote the set of all vertex sets of complete subgraphs of G . Put $\alpha(G) =: \alpha$, and let \mathcal{A} be the set of all independent vertex sets A in G with $|A| = \alpha$. Every proper induced subgraph of \bar{G} is the complement of a proper induced subgraph of G , and is

hence perfect by induction. For the perfection of \bar{G} it thus suffices to prove $\chi\bar{G} \leq \omega\bar{G}(= \alpha)$. To this end, we shall find a set $K \in \mathcal{K}$ s.t $K \cap A \neq \emptyset$ for all $A \in \mathcal{A}$; then

$$\omega(\bar{G} - K) = \alpha(G - K) + 1 = \omega(\bar{G} - K) + 1 \leq \omega(\bar{G})$$

so by the induction hypothesis

$$\chi(\bar{G}) \leq \chi(\bar{G} - K) + 1 = \omega(\bar{G} - K) + 1 \leq \omega(\bar{G})$$

as desired. Suppose there is no such K ; thus, for every $K \in \mathcal{K}$ there exists a set $A_K \in \mathcal{A}$ with $K \cap A_K = \emptyset$. Let us replace in G every vertex x by a complete graph G_x of order

$$k(x) := |\{K \in \mathcal{K} | x \in A_K\}|$$

joining all the vertices of G_x to all the vertices of G_y whenever x and y are adjacent in G . The graph G' thus obtained has vertex set $\cup_{x \in V} V(G_x)$, and two vertices $v \in G_x$ and $w \in G_y$ are adjacent in G' iff. $x = y$ or $xy \in E$. Moreover, G' can be obtained by repeated vertex expansion from the graph $G[\{x \in V | k(x) > 0\}]$. Being an induced subgraph of G , this latter graph is perfect by assumption, so G' is perfect by claim 1. In particular,

$$\chi(G') \leq \omega(G'). \quad (*1)$$

In order to obtain a contradiction (*1), we now compute in turn the actual values of $\omega(G')$ and $\chi(G')$. By construction of G' , every maximal complete subgraph of G' has the form $G'[\cup_{x \in X} G_x]$ for some $X \in \mathcal{K}$. So there exists a set $X \in \mathcal{K}$ s.t

$$\begin{aligned}
\omega(G') &= \sum_{x \in X} k(x) \\
&= |\{(x, K) : x \in X, K \in \mathcal{K}, x \in A_K\}| = 0 \\
&= \sum_{K \in \mathcal{K}} |X \cap A_K| \\
&\leq |\mathcal{K}| - 1 \cdots (*2)
\end{aligned}$$

the last inequality follows from the fact that $|X \cap A_K| \leq 1$ for all K (since A_K is independent but $G[X]$ is complete), and $|X \cap A_X| = 0$ (by the choice of A_X). On the other hand,

$$\begin{aligned}
|G'| &= \sum_{x \in V} k(x) \\
&= |\{(x, K) : x \in V, K \in \mathcal{K}, x \in A_K\}| \\
&= \sum_{K \in \mathcal{K}} |A_K| \\
&= |\mathcal{K}| \alpha
\end{aligned}$$

As $\alpha(G') \leq \alpha$ by construction of G' , this implies

$$\chi(G') \geq \frac{|G'|}{\alpha(G')} \geq \frac{G'}{\alpha} = |\mathcal{K}| \cdots (*3)$$

Putting (*2) and (*3) together we obtain

$$\chi(G') \geq |\mathcal{K}| > |\mathcal{K}| - 1 \geq \omega(G'),$$

a contradiction to (*1)

□

Graph Theory I Homework III-3 by 陳巧玲

3. Find a graph G which is C_3 -free and $\chi(G)=6$.

<sol>

I want to use Mycielski's construction to produce a graph G

which is C_3 -free and

$$\chi(G)=6.$$

Then, I want to prove a theorem:

<Thm>From a k -chromatic triangle-free graph G' , Mycielski's construction to produce a $k+1$ -chromatic triangle-free graph G .

<pf>

Let $V(G') = \{v_1, v_2, v_3, v_4, \dots, v_n\}$ and let G be a graph produced from

it by Mycielski's construction. Let u_1, u_2, \dots, u_n be the copies of

$v_1, v_2, v_3, v_4, \dots, v_n$ with w the additional vertex. Let $U = \{u_1, u_2, \dots, u_n\}$

By construction, U is an independent set in G . Hence, the other

vertices of any triangle containing u_i must belong to $V(G')$, so

let $v_{k_i}, v_{k_{i+1}} \in V(G')$, s.t. $v_{k_i} - v_{k_{i+1}} - u_i$ form a triangle. And this

is also implies $v_i - v_{k_i} - v_{k_{i+1}}$ form a triangle in G' . This triangle

can't exist in G' .

A proper k -coloring f of G' extends to a proper $k+1$ -coloring of

G , by setting $f(u_i) = f(v_i)$ and $f(w) = k+1$; hence $\chi(G) \leq \chi(G') + 1$. (1)

Now, we want to prove $\chi(G') < \chi(G)$.

Let g be the proper p -coloring of G .

By changing the names of colors, we may assume $g(w) = p$. So

$$g(u_i) \in \{1, 2, \dots, p-1\} \forall u_i \in U.$$

Then, we may use p colors to color $V(G')$. Let $A = \{v_i \mid g(v_i) = p, v_i \in V(G')\}$

We want to change color of v_j , where $v_j \in A$ s.t

$$g(v_j) \in \{1, 2, \dots, p-1\}$$

For $g(u_i) \in \{1, 2, \dots, p-1\} \forall u_i \in U$, so we change the color of v_j to $g(u_j)$.

That is because $g(v_{j_1}) = p = g(v_{j_2})$ $\{v_{j_1}, v_{j_2}\} \notin E(G')$

$$v_{j_1}, v_{j_2} \in A$$

So, we only to check $\{v'_j, v_j\} \in E(G')$ in order to conclude we

can use $p-1$ colors to color $V(G')$, where $v'_j \in G', v_j \in A$

For $\{v'_j, v_j\} \in E(G')$ where $v'_j \in G', v_j \in A$, so $\{v'_j, u_j\} \in E(G)$ which yields $g(u_j) \neq g(v'_j)$. However, $g(v_j) = g(u_j)$, so $g(v_j) \neq g(v'_j)$.

This means, the way we change color of v_j , where $v_j \in A$, can

be used.

We have shown that the modified coloring of $V(G')$ is a proper $p-1$ coloring of G' .

Hence, we can conclude $\chi(G') < \chi(G)$.

Also, by (1) $\chi(G) \leq \chi(G') + 1$, so $\chi(G') = \chi(G)$.

Due to this theorem, we can construct graph G , where

$\chi(G) = 6$, step by step.

Graph Theory I Homework III-4 by 游舜婷

4. Show that $\chi'(K_{m(n)}) = (m-1)n$ provided mn is even

Fact $K_{m(n)}$ can be decomposed into Hamiltonian cycles if

- (1). n is even, or
- (2). n is odd and m is odd

Sol: Since mn is even, we consider the following two cases

<case1> If m is odd

Since mn is even $\Rightarrow n$ is even and $\forall v \in K_{m(n)}, \deg(v) = (m-1)n$

(1). By the fact we have $K_{m(n)}$ can be decomposed into $\frac{(m-1)n}{2}$

Hamiltonian cycles C_{mn}

(2). Since n is even, C_{mn} is an even cycle $\Rightarrow \chi'(C_{mn}) = 2$

$\Rightarrow \therefore$ by (1) and (2) we have $\chi'(K_{m(n)}) = \frac{(m-1)n}{2} \times 2 = (m-1)n$

Theorem (Konig)

If G is bipartite, then $\chi'(G) = \Delta(G)$

(i.e. If $G = K_{n,n}$, then $\chi'(G) = n$)

<case2> If m is even

$\Rightarrow K_m$ has 1-factorization

$\Rightarrow K_m$ has $(m-1)$ 1-factors

\therefore 將 $K_{m(n)}$ 中每個 partite set 看成一個點 $\Rightarrow K_m$

$\Rightarrow \therefore K_m$ has $(m-1)$ 1-factors.....(1)

\therefore 將每一個 1-factors 對應到 $K_{m(n)}$ 會形成 $\frac{m}{2}$ 個 $K_{n,n}$

By Konig Theorem we have $\chi'(K_{n,n}) = n$(2)

(對每一個 1-factor 來說，我們都用這 n 個顏色去塗這 $\frac{m}{2}$

個 $K_{n,n}$ ，對不同的 1-factor 則用不同的 n 個顏色塗)

\therefore by (1) (2) we have $\chi'(K_{m(n)}) = (m-1)n$

Hence by <case1> and <case2> we have $\chi'(K_{m(n)}) = (m-1)n$

5. Find a class of infinite graphs G such that G is conformable and G is of type 2

sol : $K_{2m,2m}$, $m > 0$

claim1 : $K_{2m,2m}$ is conformable.

Defn. : If G has a vertex coloring Ψ such that

$$|\{i \mid |(\Psi^{-1}(i))| \not\equiv |V(G)| \pmod{2}\}| \leq \text{def}(G)$$

Then G is *conformable*.

By $\chi(K_{2m,2m})=2$, 不難發現任一顏色都著了 $2m$ 點。

Let Ψ 為點著色

$$|\{i \mid |(\Psi^{-1}(i))| \not\equiv |V(G)| \pmod{2}\}| = 0$$

又 $\text{def}(G)=0$ (regular graph)

$\therefore K_{2m,2m}$ is conformable

claim2 : $K_{2m,2m}$ is of type 2

由 $\chi''(G) \leq \chi(G) + \chi'(G)$ 可得知

$$\chi''(K_{2m,2m}) \leq 2 + 2m - (i)$$

接下來證 $\chi''(K_{2m,2m}) \geq 2 + 2m$

$$(i.e. \chi''(K_{2m,2m}) > 2m + 1 = \Delta(G) + 1)$$

若 $2m+1$ 色可著，則每個顏色出現 $\frac{(2m)^2+4m}{2m+1}$ 次。

By 鴿籠原理，至少存在一種顏色至少出現 $2m+1$ 次。令至少出現 $2m+1$ 次的顏色為 C_i

W.L.O.G，只討論 C_1

$\because K_{2m,2m}$ 的最大 matching 數為 $2m$ ，

\therefore 必存在一些點著色為 C_1 。

又有邊相鄰的兩點不能著同色，

\therefore 點著色為 C_1 的點在同一側。

\because 多一點著 C_1 色，可著 C_1 色的邊數就少一。

(let t 點著 C_1 色，找可著 C_1 色的邊數意同於在 $K_{2m-t,2m}$ 找 max matching 數 = $2m-t$)

$\therefore C_1$ 最多只能出現 $2m$ 次 $\rightarrow \leftarrow$

$$\text{得 } \chi''(K_{2m,2m}) \geq 2m + 2 - (ii)$$

By (i)(ii) $\chi''(K_{2m,2m}) = 2m + 2 = \Delta(G) + 2$ is type 2

By 1,2 得 $K_{2m,2m}$ 符合題意 #

Graph Theory I Homework III-6 by 李光祥

6. For $t = 10, 11, 12$, find a graph of order 30 which can be decomposed into t Hamiltonian cycles and $K_{30} - G$ contains no Hamiltonian cycle.

(sol)

Definition :

(1) $D(\{i\}, n)$ is the graph with vertex set Z_n and edge set $\{\{x, y\} \mid |x - y| \in \{i, n - i\}\}$.

(2) If $I \subseteq \{1, 2, \dots, \lfloor n/2 \rfloor\}$, then let $D(I, n) = \bigcup_{i \in I} D(\{i\}, n)$.

Theorem(Lasker and Auerbach, 1976) :

$K_{m(n)}$ has a hamiltonian decomposition if and only if $(m - 1)n$ is even.

(In fact, $K_{m(n)}$ can be decomposed into $(m - 1)n/2$ Hamiltonian cycles.)

Note :

$D(I, n)$ is connected if and only if $\gcd(I \cup \{n\}) = 1$.

1. $t = 10$

Let $G = K_{3(10)}$.

Since $(3 - 1) \times 10 = 20$ is even, by the above Theorem,

$K_{3(10)}$ can be decomposed into 10 Hamiltonian cycles.

$K_{30} - G = D(I, 30)$, where $I = \{3, 6, 9, 12, 15\}$.

Since $\gcd(I \cup \{30\}) = 3 \neq 1$, $K_{30} - G$ is disconnected.

$\Rightarrow K_{30} - G$ contains no Hamiltonian cycle.

3. $t = 12$

Let $G = K_{5(6)}$.

Since $(5 - 1) \times 6 = 24$ is even, by the above Theorem,

$K_{5(6)}$ can be decomposed into 12 Hamiltonian cycles.

$K_{30} - G = D(I, 30)$, where $I = \{5, 10, 15\}$.

Since $\gcd(I \cup \{30\}) = 5 \neq 1$, $K_{30} - G$ is disconnected.

$\Rightarrow K_{30} - G$ contains no Hamiltonian cycle.

Note :

(1) $D(I, n)$ is a Cayley graph if and only if it is connected.

(2) $D(I, n)$ is connected and 4-regular when $I = \{j - 1, j\}$,

$$2 \leq j \leq \lfloor (n-1)/2 \rfloor.$$

Theorem (Bermond, Favaron and Maheo, 1989) :

Every connected 4-regular Cayley graph on a finite abelian group has a hamiltonian decomposition.

(In fact, it can be decomposed into two Hamiltonian cycles.)

2. $t = 11$

Let $G = D(I, 30)$, where $I = \{1, 2, 4, 5, 7, 8, 9, 10, 11, 13, 14\}$.

We can partition I into $\{1, 2\}$, $\{4, 5\}$, $\{7, 8\}$, $\{9, 10\}$, $\{11\}$, $\{13, 14\}$.

Since $D(\{1, 2\}, 30)$, $D(\{4, 5\}, 30)$, $D(\{7, 8\}, 30)$, $D(\{9, 10\}, 30)$, $D(\{13, 14\}, 30)$ are connected 4-regular Cayley graphs, by the above Theorem, each of them can be decomposed into two Hamiltonian cycles.

Since $D(\{11\}, 30)$ is a connected 2-regular graph, it is a Hamiltonian cycle.

Hence G can be decomposed into 11 Hamiltonian cycles.

$K_{30} - G = D(I', 30)$, where $I' = \{3, 6, 12, 15\}$.

Since $\gcd(I' \cup \{30\}) = 3 \neq 1$, $K_{30} - G$ is disconnected.

$\Rightarrow K_{30} - G$ contains no Hamiltonian cycle. \square

Graph Theory I Homework III-7 by 李柏瑩

7. Give as many sufficient conditions for the existence of the Hamiltonian cycles in a graph as possible.

Sol:

Theorem1:

Every graph with $n > 3$ vertices and minimum degree at least $n/2$ has a Hamilton cycle. (Dirac 1952)

proof:

Let $G = (V; E)$ be a graph with $|G| = n > 3$ and $\delta(G) \geq \frac{n}{2}$.

Then G is connected

Let $P = x_0 x_1 \dots x_k$ be a longest path in G .

By the maximality of P , all the neighbors of x_0 and all the neighbors of x_k lie on P .

Hence

at least $\frac{n}{2}$ of the vertices $x_0 x_1 \dots x_{k-1}$ are adjacent to x_k , and at least

$\frac{n}{2}$ of these same $k < n$ vertices x_i are such that $x_0 x_{i+1} \in E$.

By the

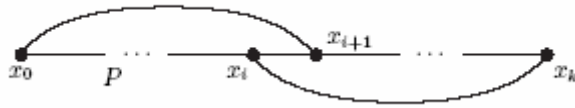
Pigeonhole principle, there is a vertex x_i that has both properties, so

we have $x_0 x_{i+1} \in E$ and $x_i x_k \in E$ for some $i < k$

We claim that the cycle $C := x_0 x_{i+1} P x_k x_i P x_0$ is a Hamilton cycle

of G . Indeed, since G is connected, C would otherwise have a neighbor

in $G-C$, which could be combined with a spanning path of C into a



path longer than P .

Theorem2:

Every graph G with $|G| \geq 3$ and $\alpha(G) \leq \kappa(G)$ has a Hamilton cycle.

Theorem3:

Every 4-connected planar graph has a Hamilton cycle.

Graph Theory I Homework III-8 by 劉宜君

8. Use Petersen's Theorem to prove that $\chi'(G) \leq 3 \left\lceil \frac{\Delta(G)}{2} \right\rceil$, when G is loopless graph.

Petersen's Theorem: Every $2k$ -regular graph has a 2-factor.

Pf. G is a loopless graph, $V(G) = \{v_1, v_2, \dots, v_m\}$

Let G' be a graph with $V(G') = \{w_1, w_2, \dots, w_m\}$, $E(G') = \{w_i w_j \mid v_i v_j \in E(G)\}$.

Let $H = G \cup G' \cup \{\Delta(G) - \deg(v_i) \text{ copies of the edge } v_i w_j\}$

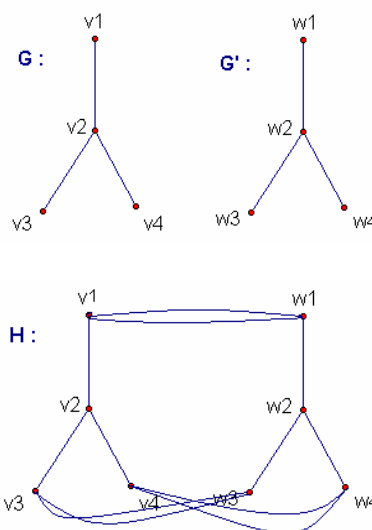
Then H be a $\Delta(G)$ -regular graph.

(i) If $\Delta(G)$ is even, then by Petersen's Theorem.

We can partition H into $\frac{\Delta(G)}{2}$ 2-factors.

Since each 2-factor is disjoint union of cycles,
and $\chi'(C_n) \leq 3$ $\left(\chi'(C_n) = \begin{cases} 2 & \text{if } n \text{ is even} \\ 3 & \text{if } n \text{ is odd} \end{cases} \right)$

$$\Rightarrow \chi'(H) \leq 3 \times \frac{\Delta(G)}{2}.$$



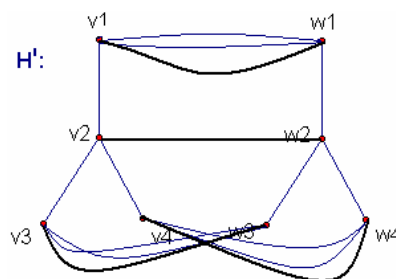
(ii) If $\Delta(G)$ is odd. Consider $H' = H \cup \{v_i w_j\}_{i=1}^m$

Then H' be a $(\Delta(G) + 1)$ -regular graph and $\Delta(G) + 1$ is even, by Petersen's Theorem.

We partition H into $\frac{\Delta(G) + 1}{2}$ 2-factors.

Since each 2-factor is disjoint union of cycles,

$$\Rightarrow \chi'(H') \leq 3 \times \frac{\Delta(G) + 1}{2}.$$



By the definition of H and H' , and (i) (ii) we can find that there exist a loopless

supergraph H_s of G , such that $\chi'(H_s) \leq 3 \left\lceil \frac{\Delta(G)}{2} \right\rceil$.

We delete those edges of $(H_s - G)$ to obtain a proper edge-coloring of G such that

$$\chi'(G) \leq 3 \left\lceil \frac{\Delta(G)}{2} \right\rceil.$$

Graph Theory I Homework III-9 by 劉士慶

9. Prove that every k -chromatic graph has at least $\binom{k}{2}$ edges. Use this to prove that

if G is the union of m complete graphs of order m , then $\chi(G) \leq 1 + m\sqrt{m-1}$.

Consider a k -coloring graph of a k -chromatic G .

I. If $e(G) < \binom{k}{2}$, this means that there exist two colors a, b such that any adjacent vertices do not require the two colors. (\Leftrightarrow the two colors would not form any edges of G). Thus the vertices with colors a and b form a single independent set, and $V(G)$ is covered by $k-1$ independent sets.

A contradiction, since a k -chromatic graph has k independent sets.

II. Since G is the union of m complete graphs of order m , which implies.

Let $\chi(G) = k$, then $\binom{k}{2} \leq m \binom{k}{2} \Rightarrow k^2 - k - m^2(m-1) \leq 0$

$$\because \sqrt{x+1} \leq \sqrt{x} + 1, \forall x \geq 0$$

$$\therefore k \leq \frac{1}{2} [1 + \sqrt{1 + 4m^2(m-1)}] \leq [1 + (1 + \sqrt{4m^2(m-1)})] = 1 + m\sqrt{m-1}$$

Hence $\chi(G) = k \leq 1 + m\sqrt{m-1}$.

Graph Theory I Homework III-10 by 曾慧棻

Graph Theory

Hui-Fen, Tseng

January 26, 2008

10. Improvement of Brooks' Theorem

(a) Given a graph G , let k_1, \dots, k_t be nonnegative integers with $\sum k_i \geq \Delta(G) - t + 1$. Prove that $V(G)$ can be partitioned into sets V_1, \dots, V_t so that for each i , the subgraph G_i induced by V_i has maximum degree at most k_i .

(b) For $4 \leq r \leq \Delta(G) + 1$, use part (a) to prove that $X(G) \leq \lceil \frac{r-1}{r}(\Delta(G) + 1) \rceil$ when G has no r -clique.

pf:

(a) By induction on t

If $t=2$, we claim the partition minimize $k_1e(G_2) + k_2e(G_1)$ satisfying $\Delta(G_1) \leq k_1$ and $\Delta(G_2) \leq k_2$

Suppose not, there exists a vertex x in G_1 s.t. $deg_G(x) > k_1$, it implies $deg_G(x) \geq k_1 + 1$. Since $k_1 + k_2 \geq \Delta(G) - 2 + 1 = \Delta(G) - 1$, i.e., $\Delta(G) \leq k_1 + k_2 + 1$, we obtain $|N_{G_2}(x)| \leq k_2$ in $V(G_2)$

Next, we move x in G_1 to G_2 to form G'_1 and G'_2 , we obtain

$$\begin{aligned} k_1e(G'_2) + k_2e(G'_1) &= k_1(e(G_2) + k_2) + k_2(e(G_1) - k_1 - 1) \\ &= k_1e(G_2) + k_1k_2 + k_2e(G_1) - k_2k_1 - k_2 \\ &= k_1e(G_2) + k_2e(G_1) - k_2 \end{aligned}$$

we find another partition $k_1e(G'_2) + k_2e(G'_1)$ is less than $k_1e(G_2) + k_2e(G_1)$, it contradicts. We have done.

If $t > 2$. Let $h = k_1 + k_2 + \dots + k_{t-1} + (t - 2)$

Suppose k_1, k_2, \dots, k_{t-1} is ok.

If $\sum_{i=1}^t k_i \geq \Delta(G) - t + 1$, it follow that

$$\begin{aligned} & k_1 + k_2 + \dots, k_{t-1} + k_t \geq \Delta(G) - t + 1 \\ \Rightarrow & h + k_t \geq \Delta(G) - t + 1 + t - 2 \\ \Rightarrow & h + k_t \geq \Delta(G) - 1 \end{aligned}$$

By $t=2$ method, we obtain $\Delta(G_h) \leq k_h, \Delta(G_t) \leq k_t$,

it follows that $\Delta(G_i) \leq k_i$ for $i=1,2,\dots,t$. We have done.

(b) $r \leq \Delta(G) + 1$ it implies, there exists t s.t.

$(\sum_{i=1}^{t-1} r) + r^* = \Delta(G) + 1$, where $2r > r^* \geq r$. Actually,
 $t = \lfloor \frac{\Delta(G)+1}{r} \rfloor$.

Let $k_1 = k_2 = \dots = k_{t-1} = r - 1. k_t \geq r - 1$ and

$\sum_{i=1}^t k_i = \Delta(G) - t + 1$, where $4 \leq r \leq \Delta(G) + 1$

$$\Rightarrow \Delta(G) - t + 1 \geq t(r - 1) = tr - t$$

$$\Rightarrow t \leq \frac{\Delta(G)+1}{r}$$

$$\Rightarrow t \leq \lfloor \frac{\Delta(G)+1}{r} \rfloor$$

By (a) $V(G)$ can partition into t classes s.t.

$\Delta(G_i) \leq k_i, i=1,\dots,t$. Since G has no r -clique, G_i has

no r -clique. $i=1,\dots,t$. Thus, by Brooks' theorem, we have

$X(G_i) \leq \Delta(G_i)$, it follows that $X(G_i) \leq k_i$ for all $i=1,\dots,t$.

Now, coloring the subgraphs with disjoint color sets, we have

$$\begin{aligned} X(G) & \leq \sum_{i=1}^t X(G_i) \leq \sum_{i=1}^t k_i = \Delta(G) + 1 - t \\ & \leq \Delta(G) + 1 - \lfloor \frac{\Delta(G) + 1}{r} \rfloor \\ & \leq \lceil \frac{r-1}{r} (\Delta(G) + 1) \rceil. \end{aligned}$$