

Graph Theory (I), Homework 3-1 by 吳政軒

Let H_i^j be the j -th Hamiltonian cycle when $k=i$

Let $V(C_p) = \{1, 2, \dots, p\}$, Define $d_i = \{(x, y) \mid d(x, y) = i, xy \in V(C_p)\}$

where $d(x, y)$ is the length of shortest path from x to y

(i) $k=2$

Case1: if p is odd

$\Rightarrow H_1^2$ 即為 d_1 構成的 cycle $\langle 1, 2, 3, \dots, p-1, p \rangle$

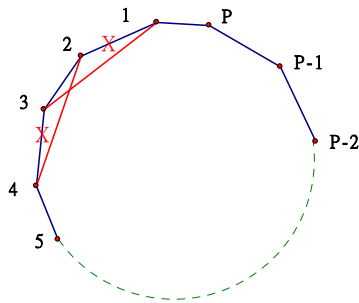
H_2^2 即為 d_2 構成的 cycle $\langle 1, 3, 5, \dots, p, 2, 4, \dots, p-1 \rangle$

Case1: if p is even

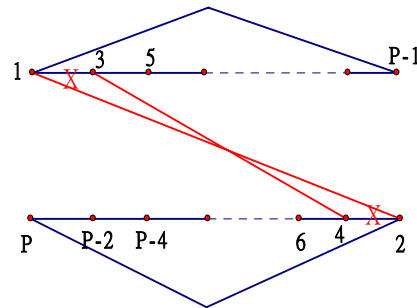
$\Rightarrow H_1^2$ 即為 $d_1 - \{(1, 2), (3, 4)\} + \{(1, 3), (2, 4)\}$ 構成的 cycle

H_2^2 即為 $d_2 - \{(1, 3), (2, 4)\} + \{(1, 2), (3, 4)\}$ 構成的 cycle

H_1^2



H_2^2



(ii) $k=3$

Case1: if $(p,3)=1$, i.e. $p \equiv 1$ or $2 \pmod{3}$

$\Rightarrow H_1^3 = H_1^2$, $H_2^3 = H_2^2$, H_3^3 即為 d_3 構成的 cycle

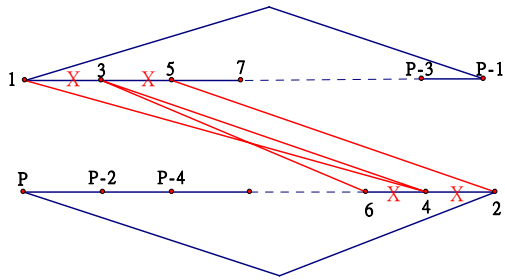
Case2: if $p \equiv 0 \pmod{2}$ and $p \equiv 0 \pmod{3}$

$\Rightarrow H_1^3 = H_1^2$,

H_2^3 即為 $d_2 - \{(1,3), (2,4), (3,5), (4,6)\} + \{(1,4), (2,5), (3,6), (3,4)\}$ 構成的 cycle

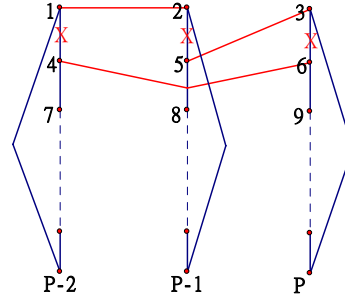
H_3^3 即為 $d_3 - \{(1,4), (2,5), (3,6)\} + \{(3,5), (4,6), (1,2)\}$ 構成的 cycle

H_2^3



$\langle 1,4,3,6,8,\dots,p,2,5,7,\dots,p-1 \rangle$

H_3^3



Case3: if $p \equiv 0 \pmod{3}$ and $p \equiv 1 \pmod{2}$

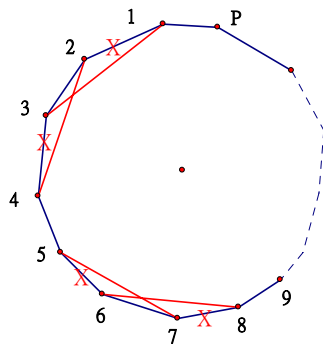
$\Rightarrow H_1^3$ 即為 $d_1 - \{(1,2), (3,4), (5,6), (7,8)\} + \{(1,3), (2,4), (5,7), (6,8)\}$ 構成的 cycle

H_2^3 即為

$d_2 - \{(1,3), (2,4), (3,5), (4,6), (5,7), (6,8)\} + \{(1,4), (2,5), (3,6), (3,4), (5,6), (7,8)\}$ 構成的 cycle

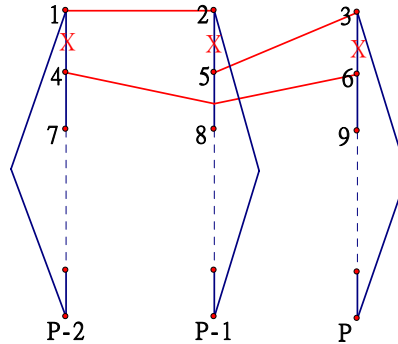
H_3^3 即為 $d_3 - \{(1,4), (2,5), (3,6)\} + \{(3,5), (4,6), (1,2)\}$ 構成的 cycle

H_1^3

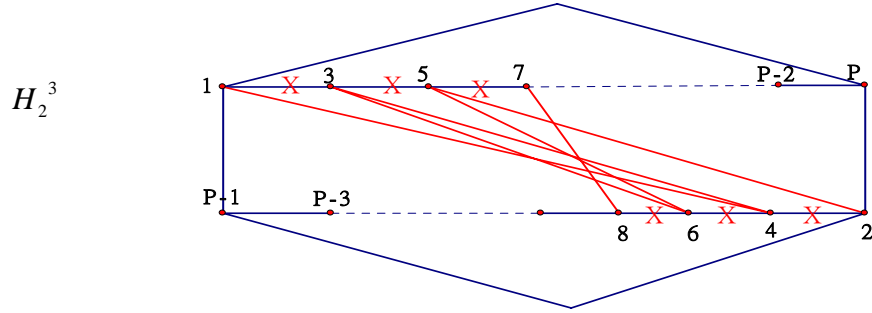


$\langle 1,3,2,4,5,7,6,8,9,\dots,p \rangle$

H_3^3



$\langle 1,2,p-1,p-4,\dots,5,3,p,p-3,\dots,6,4,7,\dots,p-2 \rangle$



$$\langle 1,4,3,6,5,2,p,p-2,\dots,7,8,10,\dots,p-3,p-1 \rangle$$

(iii) $k=4$

Case1: $p \equiv 1, 3 \pmod{4}$

$\Rightarrow H_1^4 = H_1^3, H_2^4 = H_2^3, H_3^4 = H_3^3, H_4^4$ 爲 d_4 構成的 cycle

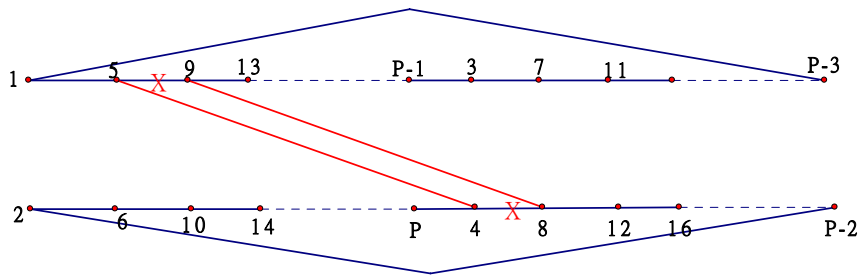
Case2: $p \equiv 2 \pmod{4}$

$\Rightarrow H_1^4$ 爲 $H_1^3 - \{(4,5), (8,9)\} + \{(4,8), (5,9)\}$ 構成的 cycle

$$H_2^4 = H_2^3, H_3^4 = H_3^3$$

H_4^4 爲 $d_4 - \{(4,8), (5,9)\} + \{(4,5), (8,9)\}$ 構成的 cycle

H_4^4



Case3: $p \equiv 0 \pmod{4}$

$\Rightarrow H_1^4$ 爲 $H_1^3 - \{(4,5), (6,7), (8,9), (9,10)\} + \{(4,8), (5,9), (6,10), (7,9)\}$

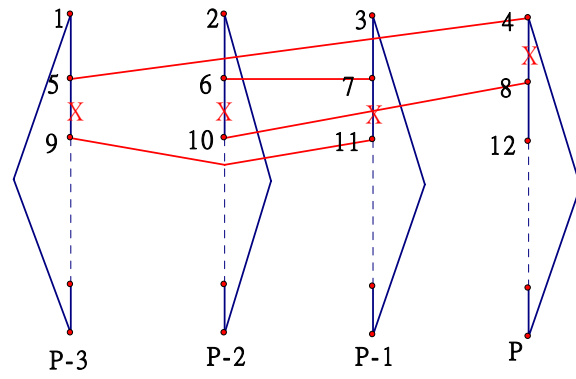
構成的 cycle

H_2^4 爲 $H_2^3 - \{(7,9), (8,10), (9,11)\} + \{(7,11), (8,9), (9,10)\}$ 構成的 cycle

H_4^4 爲 $d_4 - \{(4,8), (5,9), (6,10), (7,11)\} + \{(4,5), (6,7), (8,10), (9,11)\}$ 構成

的 cycle and $H_3^4 = H_3^3$

H_4^4



$\langle 1, 5, 4, p, p-4, \dots, 12, 8, 10, 14, \dots, p-2, 2, 6, 7, 3, p-1, p-4, \dots, 11, 9, 13, \dots, p-3 \rangle$

Graph Theory (I), Homework 3-2 by 曾雅榕

2. Let D be a directed graph such that for each

$$v \in V(D), \deg_D^+(v) = \deg_D^-(v) = t \geq 1.$$

Prove that D can be decomposed into t directed 1-factors.(with indegree 1 and outdegree 1)

Proof:

Assume D is connected and $V(D) = \{v_1, v_2, v_3, \dots, v_n\}$

Otherwise, consider each component of D

$$\text{Since } v \in V(D), \deg_D^+(v) = \deg_D^-(v) = t \geq 1$$

$\Rightarrow D$ is strong connected.

$\Rightarrow D$ has a directed E.C

Called $C: x_1 \rightarrow x_2 \rightarrow x_3 \rightarrow \dots \rightarrow x_q$ where $q = \|D\|$

Define a bipartite graph H with two partite set A, B

Where $A = \{a_1, a_2, a_3, \dots, a_n\}, B = \{b_1, b_2, b_3, \dots, b_n\}$

$$a_{x_i} \leftrightarrow b_{x_j} \text{ in } H \Leftrightarrow x_i \rightarrow x_j \text{ in } C$$

$\Rightarrow H$ is a t -regular bipartite graph

By König's Theorem

$\Rightarrow H$ has an 1-factor M

$\Rightarrow M$ in H corresponds to a directed 1-factor in D

$\Rightarrow D$ can be decomposed into t 1-factors.

Graph Theory (I), Homework 3-3 by 陳柏澍

3. Prove that G is a perfect graph if and only if \overline{G} is a perfect graph.

Def. We say that G' is obtained from G by expanding the vertex $x \in V(G)$

$$\text{if and only if } \begin{cases} V(G') = V(G) \cup \{x'\} \\ E(G') = E(G) \cup \{x'v \mid v \in N_G[x]\} \end{cases}.$$

Lemma: If G is perfect graph, then G' obtained from expanding a vertex $x \in V(G)$ from G is again a perfect graph.

pf of Lemma:

Prove by induction on $|G|$:

Expanding the vertex of K_1 yields K_2 which is clearly perfect.

Suppose G is a nontrivial perfect graph and G' is obtained from G by expanding the vertex $x \in V(G)$.

Claim: $\forall H \leq G', \chi(H) = \omega(H)$.

Case1. $H < G'$

$$\text{i. } |V(H) \cap \{x, x'\}| \leq 1$$

H is isomorphic to an induced subgraph of G .

Such G is perfect, then we are done.

$$\text{ii. } \{x, x'\} \subseteq V(H)$$

H is obtained from $H' := H - x'$ where $|H'| < |G|$

by expanding a vertex $x \in V(H')$. Then, by

induction hypothesis, done.

Case2. $H = G'$

Claim: $\chi(G') \leq \omega(G')$

Let $\omega(G) := \omega$.

Observe that $\omega(G') \in \{\omega, \omega + 1\}$.

Case2.1. $\omega(G') = \omega + 1$

$$\chi(G') \leq \chi(G) + 1 = \omega(G) + 1 = \omega + 1 = \omega(G')$$

Clearly, $\chi(G') \geq \omega(G')$.

$$\therefore \chi(G') = \omega(G').$$

Case 2.2. $\omega(G') = \omega$

Note, x lies in no $K_\omega \subseteq G$, otherwise $\omega(G') = \omega + 1$.

Coloring G with ω colors.

Let X be a color class containing x .

Then, $\begin{cases} |X \cap V(K_\omega)| = 1 \\ x \notin X \cap V(K_\omega) \end{cases}$ for any $K_\omega \subseteq G$.

(Clear, $|X \cap V(K_\omega)| \leq 1$. If $X \cap V(K_\omega) = \emptyset$, then ω colors will not be enough to color G .)

Let $\bar{G} = G - (X \setminus \{x\})$.

Thus, $\chi(\bar{G}) = \omega - 1$.

$\therefore V(G') \setminus V(\bar{G}) = (X \setminus \{x\}) \cup \{x'\}$ is an independent set.

$$\therefore \chi(G') \leq (\omega - 1) + 1 = \omega.$$

Thus, $\chi(G') = \omega(G')$. ($\therefore \chi(G') \geq \omega(G')$)

pf:

It suffices to show that if G is perfect then \bar{G} is also perfect.

Applying induction on $|G|$:

$|G| = 1$ is clear.

Suppose it holds for $|G| \geq 2$.

It suffices to show that $\chi(\bar{G}) \leq \omega(\bar{G})$.

(\therefore Every proper induced subgraph of \bar{G} is the complement of a proper induced subgraph of G .
 \therefore By induction hypothesis, every proper induced subgraph of \bar{G} is perfect.)

Let $K = \{S \subseteq V(G) \mid \langle S \rangle_G = \text{maximal } K_t \text{ for some } t \leq \omega(G)\}$

$A = \{A \subseteq V(G) \mid A \text{ is an independent set s.t. } |A| = \alpha(G)\}$.

Claim: $\exists K \in K$ s.t. $K \cap A = \emptyset$ for all $A \in A$.

pf of claim:

Suppose not.

i.e. $\forall K \in \mathbf{K} \exists A_K \in \mathbf{A}$ s.t. $K \cap A_K = \emptyset$.

Constructing G' by replacing each vertex x of G by a complete graph G_x of order $k(x) := |\{K \in \mathbf{K} | x \in A_K\}|$ and joining all the vertices of G_x to all the vertices of G_y if x and y are adjacent in G .

Observe that G' can be obtained by repeating vertex expansion from the graph $G[\{x \in V(G) | k(x) > 0\}]$ which is perfect since G is perfect.

By the Lemma, G' is perfect.

$\therefore \chi(G') = \omega(G')$. $\therefore \chi(G') \leq \omega(G')$.

Observe $\omega(G') = \sum_{x \in X} k(x)$ where $X \in \mathbf{K}$

$$= |\{(x, K) | x \in X, K \in \mathbf{K}, x \in A_K\}|$$

$$= \sum_{K \in \mathbf{K}} |X \cap A_K| \leq |\mathbf{K}| - 1.$$

$$|G'| = \sum_{x \in V(G')} k(x)$$

$$= |\{(x, K) | x \in V(G), K \in \mathbf{K}, x \in A_K\}|$$

$$= \sum_{K \in \mathbf{K}} |A_K|$$

$$= |\mathbf{K}| \cdot \alpha(G).$$

$$\therefore \alpha(G') \leq \alpha(G).$$

$$\therefore \chi(G') \geq \frac{|G'|}{\alpha(G')} \geq \frac{|G'|}{\alpha(G)} = |\mathbf{K}|.$$

Thus, $\chi(G') \geq |\mathbf{K}| > |\mathbf{K}| - 1 \geq \omega(G')$. A contradiction.

Thus, $\omega(\overline{G} - K) = \alpha(G - K) < \alpha(G) = \omega(\overline{G})$.

$\therefore \chi(\overline{G}) \leq \chi(\overline{G} - K) + 1 = \omega(\overline{G} - K) + 1 \leq \omega(\overline{G})$.

Graph Theory (I), Homework 3-4 by 張澍仁

1° Let $k = \chi'(G)$

If $k \leq \frac{3}{2}\Delta(G)$, then done.

Assume that $k > \frac{3}{2}\Delta(G)$.

2° $\because \chi'(G) = k$

$\therefore \exists G'$ be a minimal subgraph of G

s.t. $\chi'(G') = k$ and $\chi'(G' - e) = k - 1 \quad \forall e \in E(G')$.

3° Let $x, y \in V(G')$ s.t. there are $\mu(G')$ edges between x, y .

Let e be an edge of them.

4° Consider the graph: $G' - e$

We found that:

1) $\chi'(G' - e) = k - 1$

2) all x, y 所連的邊(in $G' - e$)一定要用完 $k-1$ 色.

(else by 1), 接回 e 並著未用到的顏色 $\Rightarrow \chi'(G') = k - 1 \rightarrow \leftarrow$)

3) $\because d_{G'-e}(x) = d_{G'}(x) - 1 \leq \Delta(G) - 1$.

$\therefore x$ 最多用到 $\Delta(G) - 1$ 色 (其中與 y 共用 $\mu(G') - 1$ 色).

$\Rightarrow x$ 至少還有 $(k - 1) - (\Delta(G) - 1) = k - \Delta(G)$ 色未用到.

\Rightarrow by 2), y 一定要用到這些顏色(這些顏色不在 x, y 之間).

同理, x 亦至少用到 $k - \Delta(G)$ 色, 且這些顏色不在 x, y 之間.

$\Rightarrow k - 1 \geq (\mu(G') - 1) + 2(k - \Delta(G)). \quad \dots(*)$

4° By Vizing Theorem

$$\Rightarrow k = \chi'(G) = \chi'(G') \leq \Delta(G') + \mu(G') \leq \Delta(G) + \mu(G')$$

$$\Rightarrow \mu(G') \geq k - \Delta(G)$$

$$\therefore (*) \Rightarrow k - 1 \geq (k - \Delta(G) - 1) + 2(k - \Delta(G))$$

$$\Rightarrow k \leq \frac{3}{2}\Delta(G). \quad \rightarrow\leftarrow$$

□

Graph Theory (I), Homework 3-5 by 陳宏賓

6. Prove that if $\|G\| = \binom{n+1}{2}$ and $\Delta(G) \leq \lfloor n/2 \rfloor$, then G can be decomposed into n subgraphs G_1, G_2, \dots, G_n such that G_i is induced by a matching with i edges for all i .

Lemma 1: Every k -edge colorable graph has an equitable k -edge coloring.

Proof.

Case 1: n is even.

By Vizing's Theorem and Lemma 1, G has an equitable $(\frac{n}{2}+1)$ -edge coloring with color classes $A_1, A_2, \dots, A_{\frac{n}{2}+1}$ s.t. $|A_i| = n-1$ for $i = 1, \dots, \frac{n}{2}$ and $|A_{\frac{n}{2}+1}| = n$. Choose $B_i \subseteq A_i$ with $|B_i| = i$ for $i = 1, \dots, \frac{n}{2} - 1$. Let $C_i = A_i \setminus B_i$. Then $\{B_1, B_2, \dots, B_{\frac{n}{2}-1}, C_{\frac{n}{2}-1}, \dots, C_2, C_1, A_{\frac{n}{2}}, A_{\frac{n}{2}+1}\}$ is a matching decomposition as desired.

Case 2: n is odd.

The proof is similar to that of case 1. ■

Graph Theory (I), Homework 3-6 by 張雁婷

6. Let D be an n -regular digraph of order $2n + 1$, $n \geq 1$. Prove it or disprove that D has a directed Hamiltonian cycle. (D is n -regular if $\forall v \in V(D)$, $d^+(v) = d^-(v) = n$.)

Pf.

Note. The underlying graph of D is K_{2n+1} , and hence D is a tournament of order $2n + 1$.

Claim: D has a directed Hamiltonian cycle.

Pf.

Suppose not, let C be the maximum cycle in D , and $|V(C)| \leq 2n$. Let u be a vertex not in C . Since D is a tournament, there is a directed edge between u and each vertex in C .

Case (i)

If there are two edges between u and C of different directions, we can find two adjoining vertices x and y in C such that the directions from u to x and from u to y are different. Hence we have a cycle larger than C , a contradiction.

Case (ii) the directions from u to each vertex in C are the same.

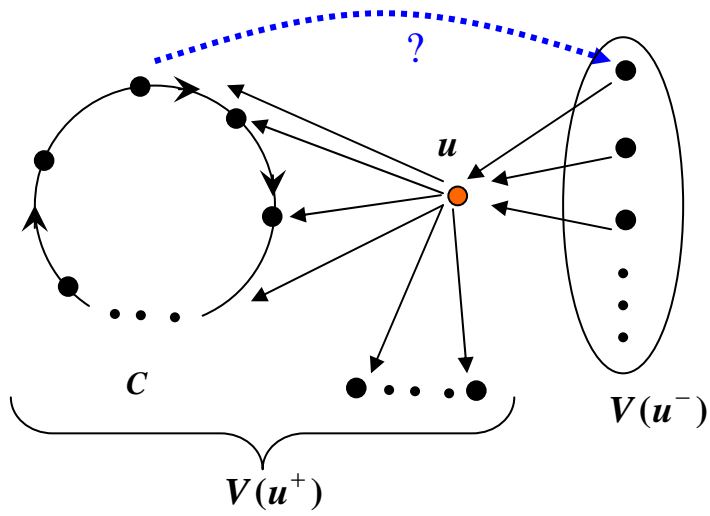
$$\therefore d^+(u) = d^-(u) = n \quad \therefore |V(C)| \leq n$$

Case (ii)-1 Suppose the direction between u and C are all $u \rightarrow C$.

$$\text{Let } V(u^-) = \{z \in V(D) \mid z \rightarrow u \text{ in } D\} \text{ and } V(u^+) = \{z \in V(D) \mid z \leftarrow u \text{ in } D\}.$$

Note that $|V(u^-)| = |V(u^+)|$. For each vertex $z \in V(u^-)$, z is also adjacent to each vertex in C , since D is a tournament. If there is a directed edge from C to z , then we have a larger cycle through u , a contradiction. Hence the direction between C and z are all $z \rightarrow C$. And $\therefore |V(u^-) + \{u\}| = n + 1$ \therefore for each

vertex $x \in C$, $\deg^-(x) \geq n + 1$, it contradicts to that D is n -regular. Therefore, D has a directed Hamiltonian cycle.



Case (ii)-2 Similar discussion for the directions between u and C are all $C \rightarrow u$.

Graph Theory (I), Homework 3-7 by 羅元勳

7. Let $n = k(2l + 1)$. Construct a non-Hamiltonian complete k -partite graph with n vertices and minimum degree $\frac{n k - 1}{2} \frac{2l}{2l + 1}$.

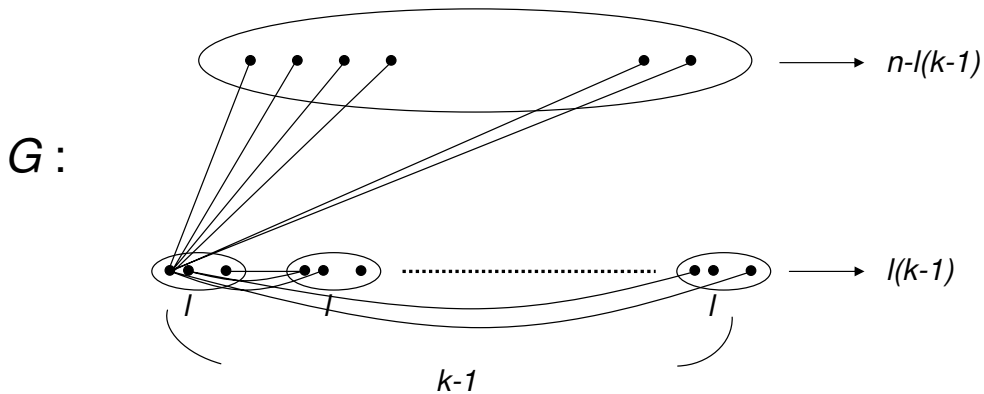
Sol.

$$\text{Minimum degree} = \frac{k(2l + 1) k - 1}{2} \frac{2l}{2l + 1} = l(k - 1).$$

Define $G = K_{n-l(k-1), l, l, \dots, l}$, which is a complete k -partite graph.

Claim : G has no Hamiltonian cycle.

$$\begin{aligned} \text{Proof. } n - l(k - 1) &= 2kl + k - kl + l \\ &= kl + k + l = l(k + 1) + k > l(k - 1) \end{aligned}$$



Thus, the vertex number of the first partite set is larger than the sum of the others, then it's easy to see that G has no Hamiltonian cycles.

Graph Theory (I), Homework 3-8 by 黃皓文

8. Dirac proved that every 2-connected simple graph G has a cycle of length at least $\min\{n(G), 2\delta(G)\}$. Use this to prove that every $2k$ -regular graph with $4k+1$ vertices is Hamiltonian.

Proof: Let G be a $2k$ -regular graph with $4k+1$ vertices

Claim: G is 2-connected

Pf of claim: Suppose not. Then there is a vertex $u \in V(G)$ s.t. $G-u$ is

not connected. Let G_1, \dots, G_t be components of $G-u$. Since G is

$2k$ -regular. $|G_i| \geq 2k$ for each i and $|G| = \sum_{i=1}^t |G_i| + 1 \geq 2tk + 1$

$\Rightarrow 4k + 1 \geq 2tk + 1 \Rightarrow t \leq 2$

Hence $t=2$ and $|G_1| = |G_2| = 2k$.

As $\deg_G(u) = 2k$, there is one vertex $v_1 \in V(G_1)$ s.t. $\deg_G(v_1) = 2k$

$\Rightarrow |G_1| \geq 2k + 1$, a contradiction.

Thus, G is 2-connected

By the result of Dirac, G has a cycle of length at least $\min\{4k+1, 4k\} = 4k$

Let C be a longest cycle in G .

If $|C| = 4k+1$, then C is a Hamiltonian cycle of G .

If $|C| = 4k$, say $C: x_1 - y_1 - x_2 - y_2 - \dots - x_{2k} - y_{2k}$, there is only one vertex

$z \in V(G-C)$. By maximality of C and $\deg_G(z) = 2k$, we see either

$N_G(z) = \{x_1, \dots, x_{2k}\}$ or $N_G(z) = \{y_1, \dots, y_{2k}\}$. W.L.O.G. assume

$N_G(z) = \{x_1, \dots, x_{2k}\}$

Claim: $y_i \sim y_j$ for some i, j

Pf of claim: Suppose not. i.e. $\forall i \neq j \in \{1, \dots, 2k\}$ y_i is not adjacent to y_j .

Then $N_G(y_i) \subseteq \{x_1, \dots, x_{2k}\}$ for each i .

As $\deg_G(y_i) = 2k$, $N_G(y_i) = \{x_1, \dots, x_{2k}\}$ for each i .

$\Rightarrow \{z, y_1, y_2, \dots, y_{2k}\} \subseteq N_G(x_1)$, a contradiction.

By claim, there is a longer cycle

$z - x_j - y_{j-1} - x_{i-1} - \dots - y_i - y_j - x_{j+1} - \dots - x_i - z$, a contradiction.

Graph Theory (I), Homework 3-9 by 黃志文

9.

- (a) $G \square K_2 = G \cup G' \cup \{v_1v'_1, v_2v'_2, v_3v'_3, \dots, v_nv'_n\}$, where G' is a copy of G and $V(G) = \{v_1, v_2, v_3, \dots, v_n\}$ and $V(G') = \{v'_1, v'_2, v'_3, \dots, v'_n\}$.

Case 1: $\chi'(G) = \Delta(G)$

Let $v_i v'_j$ be colored the same color of $v_i v_j$ and $v_i v'_i$ be colored a new color.

Then $\chi'(G \square K_2) = \Delta(G) + 1 = \Delta(G \square K_2)$.

Case 2: $\chi'(G) = \Delta(G) + 1$

Let $v_i v'_j$ be colored the same color of $v_i v_j$. Since $\chi'(G) = \Delta(G) + 1$, there is a color doesn't appear near v_i . Thus color $v_i v'_i$ with the color.

Then $\chi'(G \square K_2) = \Delta(G) + 1 = \Delta(G \square K_2)$.

- (b) It is obviously $V((G_1 \cup G_2) \square (H_1 \cup H_2)) = V((G_1 \square H_2) \cup (G_2 \square H_1))$ and let the vertex set is $\{(v, w) \mid v \in V, w \in W\}$.

1° $\forall e = \{(v, w), (v', w')\} \in E((G_1 \cup G_2) \square (H_1 \cup H_2))$.

(i) $v = v'$ and $(ww' \in E(H_1) \text{ or } ww' \in E(H_2))$.

$\Rightarrow (v = v' \text{ and } ww' \in E(H_1)) \text{ or } (v = v' \text{ and } ww' \in E(H_2))$.

$\Rightarrow e \in (G_1 \square H_2) \cup (G_2 \square H_1)$.

(ii) $w = w'$ and $(vv' \in E(G_1) \text{ or } vv' \in E(G_2))$.

$\Rightarrow (w = w' \text{ and } vv' \in E(G_1)) \text{ or } (w = w' \text{ and } vv' \in E(G_2))$.

$\Rightarrow e \in (G_1 \square H_2) \cup (G_2 \square H_1)$.

Hence $E((G_1 \cup G_2) \square (H_1 \cup H_2)) \subseteq E((G_1 \square H_2) \cup (G_2 \square H_1))$.

2° $\forall e \in E(G_1 \square H_2) \Rightarrow e \in E((G_1 \cup G_2) \square (H_1 \cup H_2))$.

$\forall e \in E(G_2 \square H_1) \Rightarrow e \in E((G_1 \cup G_2) \square (H_1 \cup H_2))$.

Hence $E((G_1 \square H_2) \cup (G_2 \square H_1)) \subseteq E((G_1 \cup G_2) \square (H_1 \cup H_2))$

So, $E((G_1 \square H_2) \cup (G_2 \square H_1)) = E((G_1 \cup G_2) \square (H_1 \cup H_2))$.

(c)

1° $\forall (v, w) \in V(G \square H)$

$\{(v, w), (v, w')\} \in E(G \square H)$ iff $w' \in N_H(w)$.

$\{(v, w), (v', w)\} \in E(G \square H)$ iff $v' \in N_G(v)$.

Hence $\deg((v, w)) = \deg(v) + \deg(w)$.

Thus $\Delta(G \square H) = \Delta(G) + \Delta(H)$.

2° Since both G and H are 1-factors, let $|G| = 2n$, $|H| = 2m$ and $G = G' \cup nK_2$, $H = H' \cup mK_2$.

Then $G \square H = (G' \cup nK_2) \square (H' \cup mK_2)$

$= (G' \square mK_2) \cup (H' \square nK_2)$ by (b)

$= (m(G' \square K_2)) \cup (n(H' \square K_2))$

And $\chi'(G \square K_2) = \Delta(G \square K_2) = \Delta(G)$

$\chi'(H \square K_2) = \Delta(H \square K_2) = \Delta(H)$ by (a).

Then $\chi'(G \square H) = \chi'((m(G \square K_2)) \cup (n(H \square K_2)))$
 $\leq \chi'(m(G \square K_2)) + \chi'(n(H \square K_2))$
 $= \chi'(G \square K_2) + \chi'(H \square K_2)$
 $= \Delta(G) + \Delta(H)$
 $= \Delta(G \square H)$
 $\leq \chi'(G \square H)$.

Thus $\chi'(G \square H) = \Delta(G \square H)$.

Graph Theory (I), Homework 3-10 by 張惠蘭

10. Prove that if G is a multigraph with multiplicity μ , then $\chi'(G) \leq \Delta(G) + \mu$.

$$\Delta = \Delta(G), \quad \mu = \mu(G)$$

1. Color G with $\Delta + \mu$ colors as many edges as possible.

Denote such coloring as f .

2. Consider some edges left uncolored.

Let $B(v)$ be the set of colors missing at v under f .

It is clear that $|B(v)| \geq \mu$ for any v .

3. Let uv be an edge which is not colored.

Claim: We can extend the coloring to uv by recoloring some edges colored under f .

Proof.

(We say that “we can have the extension” if we can extend the coloring to uv)

Let $v = v_0$.

Let $a_0 \in B(u)$, $a_1 \in B(v_0)$ and delete a_1 from $B(v_0)$.

Assume $a_1 \notin B(u)$, i.e., $\exists v_1$ s.t. $f(uv_1) = a_1$, otherwise, we can have the extension.

Let $a_2 \in B(v_1)$ and delete a_2 from $B(v_1)$.

Similarity, assume $\exists v_2$ s.t. $f(uv_2) = a_2$.

Let $a_3 \in B(v_2)$ and delete a_3 from $B(v_2)$.

Similarity, assume $\exists v_3$ s.t. $f(uv_3) = a_3$.

Notice that v_3 and v_1 are not necessarily distinct. (see fig 1)

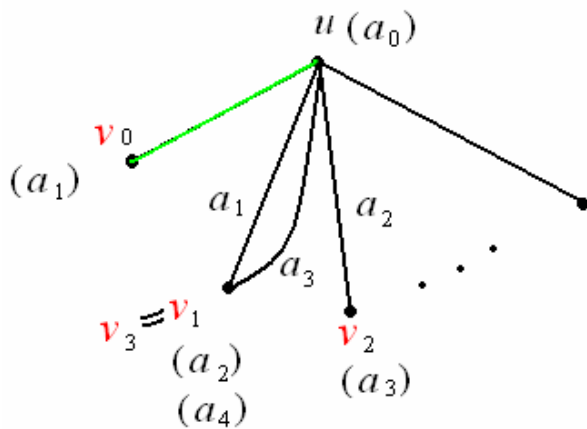


Figure 1.

But, we can always take a color from $B(v)$ in this process, since $|B(v)| \geq \mu$.

(This is different from simple graph)

Now, continue this process and then meet one of the following two cases:

- [1] Assume there exists l s.t. $l+1$ is the smallest index s.t. $a_l \in B(v_l)$ and a_l is deleted from some $B(v_{k-1})$, i.e., $a_{l-1} = a_k$.

Since a_k is deleted from $B(v_{k-1})$ before we find a_{l-1} in $B(v_l)$, $v_{k-1} \neq v_l$.

Let P be the maximal alternating path of edges colored a_0 and a_k that begins at v_l along color a_0 . Such path is well-defined, since each vertex has at most one incident edge in each color.

假設我們可以把 uv_l 的顏色從 a_l 改成 a_0 properly, 且不改變 $B(v_i)$ for any $i < t$, 則把 uv_i 塗成 a_{i+1} for $i = 0, \dots, t-1$ 仍可維持 proper coloring, 因為原本 $a_{i+1} \in B(v_i)$.----- \otimes

Consider the following cases:

- (1) $V(P) \cap \{v_0, \dots, v_{l-1}\} = \emptyset$.

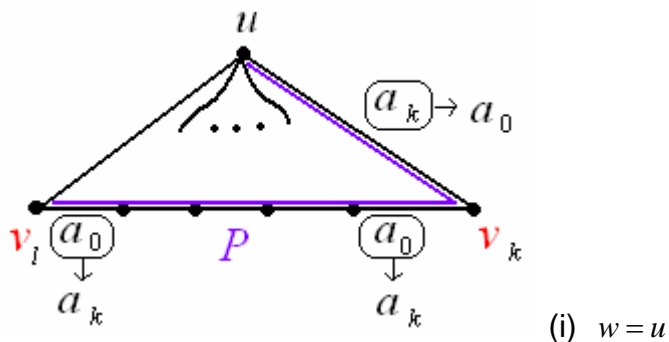
Interchange colors a_0 and a_k on P and interchange a_0 and a_l on uv_l .

It is clear that such coloring makes uv_l satisfying the condition of \otimes .

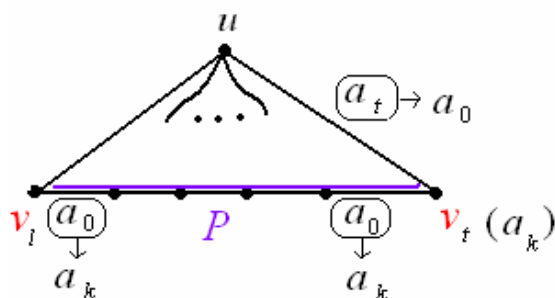
- (2) $V(P) \cap \{v_0, \dots, v_{l-1}\} \neq \emptyset$

Let w be another endpoint of P .

Then there are only three sub-cases:



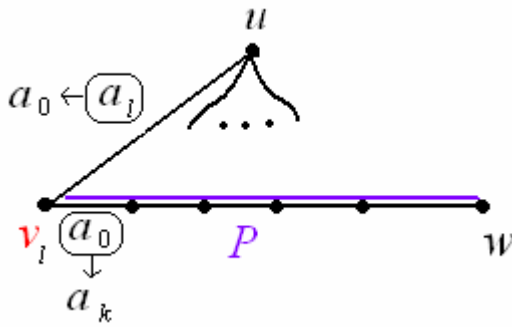
Interchange colors a_0 and a_k on P . Such coloring makes uv_k satisfying the condition of \otimes .



- (ii) $w = v_t$ for some t and $a_k \in B(v_t)$.

Interchange colors a_0 and a_k on P and change the color of an edge

colored a_t with u and v_t as its endpoints to a_0 . Such coloring makes uv_t satisfying the condition of \otimes .



(iii) $w \notin \{u, v_0, \dots, v_l\}$

Interchange colors a_0 and a_k on P and change the color of an edge colored a_l with u and v_l as its endpoints to a_0 .

Such coloring makes uv_l satisfying the condition of \otimes .

Thus if such l exist, we can always extend the coloring to uv

[2] Assume that such l doesn't exist.

Since u has at most Δ incident edges, the process can stop at some v_t s.t.

$a_{t+1} \in B(u)$.

Interchange colors a_t and a_{t+1} an edge with endpoints as u and v_t .

Such coloring makes uv_t satisfying the condition of \otimes .

Hence we can also extend the coloring to uv .